

## Digital Signatures

- ❑ Used to proof authenticity
- ❑ Alternative use of a public key cryptosystem
  - e.g., RSA
- ❑ Exchange encryption/decryption steps
  - Often used on a message digest
- ❑ More details to come ...

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## Domain Name System (DNS)

- ❑ Map
  - Human memorable hostname  
e.g.: www.net.in.tum.de
  - Machine usable addresses  
e.g.: 131.159.15.242
- ❑ DNS system
  - Application vital to the Internet
  - Transparent to the user
  - Distributed, hierarchical, redundant
  - Uses caching

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## Need for "secure" DNS

- ❑ DNS: no mechanism for authentication
- ❑ 1990's real world attacks:
  - Show how easy misuse of DNS is
  - Attacks:
    - Spoofing
    - Sniffing + answer injection
    - Guessing and predicting query Ids
    - Cache poisoning
  - Bellovin Usenix95:  
"Using the Domain Name System for System BreakIns"
  - Vixie: Usenix95:  
"DNS and BIND Security Issues"

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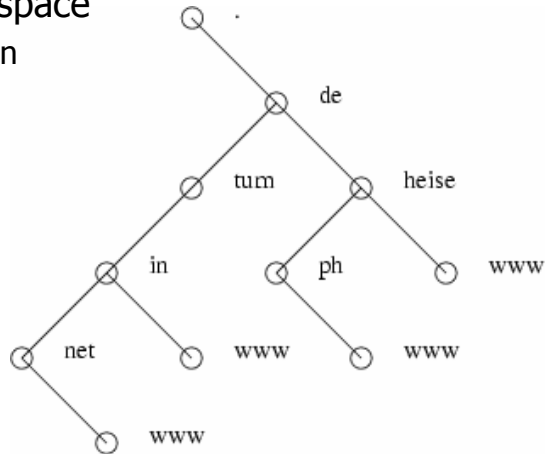
## DNSSEC: "secure" version of DNS

- ❑ Developed over the last 10 years
- ❑ Goals:
  - full backward compatibility
  - Data integrity
  - Data authenticity
- ❑ Problems:
  - Overhead???
    - Bandwidth
    - Resource consumption: Memory, CPU
  - Deployment issues
    - Key management (PKI)
    - Configuration

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## DNS

- ❑ Distributed database
- ❑ Hierarchical name space
  - Based on delegation of responsibility



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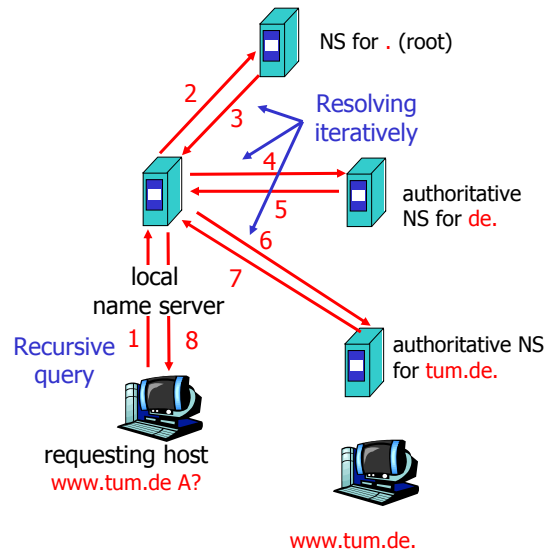
## DNS protocol

- ❑ Query ↔ Response (1-to-1 relationship)
- ❑ UDP (TCP as fallback)
- ❑ DNS message
  - Four sections
    - Query
    - Answer
    - Authority
    - Additional
- ❑ Resolution process
  - Iterative
  - Recursive

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## DNS example

- ❑ Iterative queries
- ❑ Intermediate name servers
  - Delegated zones
  - Query referral
- ❑ Information
  - Resource records (RR)
  - Can be cached
- ❑ Query types
  - recursive
  - iterative



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## DNSSEC (Delegation Signer) Basics

- ❑ Idea
  - Authentication chain: root zone to resource record (RR)
  - Parent zone guarantees child zone key
- ❑ Realization
  - Four new record types
    - **RRSIG** signatures for resource records (RRs)
    - **DNSKEY** public key for the zone (e.g., RSA)
    - **DS** digest of the child-zone's DNSKEY (at delegation points in parent zone child's name equals DS RR's name)
    - **NSEC** needed for certifying non existence
  - Upward compatibility
    - Islands of security

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## Resolving with DNS/DNSSEC

- ❑ Resolving A [www.tum.de](http://www.tum.de): empty resolver cache

DNS

Zone:

NS:

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## Resolving with DNS

- ❑ Resolving A [www.tum.de](http://www.tum.de): empty resolver cache

DNS

Zone:

NS:            .

                  de.

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## Resolving with DNS

- ❑ Resolving A [www.tum.de](http://www.tum.de): empty resolver cache

DNS

Zone: . de.

NS: de. tum.de.

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## Resolving with DNS

- ❑ Resolving A [www.tum.de](http://www.tum.de): empty resolver cache

DNS

Zone: . de. tum.de.

NS: de. tum.de.

A: [www.tum.de](http://www.tum.de)

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## Resolving with DNSSEC

- ❑ Resolving A **www.tum.de.**: empty resolver cache

### DNS/DNSSEC

Zone: .  
NS: de.  
DS: de.  
RRSIG DS: de.

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## Resolving with DNSSEC

- ❑ Resolving A **www.tum.de.**: empty resolver cache

### DNS/DNSSEC

Zone: . de.  
NS: de. tum.de.  
DS: de. tum.de.  
RRSIG DS: de. tum.de.

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## Resolving with DNSSEC

- ❑ Resolving A **www.tum.de.**: empty resolver cache

### DNS/DNSSEC

Zone:	.	de.	tum.de.
NS:	de.	tum.de.	
DS:	de.	tum.de.	
RRSIG DS:	de.	tum.de.	
RRSIG A:			www.tum.de.
DNSKEY:			tum.de.
A:			www.tum.de

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## Resolving with DNSSEC

- ❑ Resolving A **www.tum.de.**: empty resolver cache

### DNS/DNSSEC

Zone:	.	de.	tum.de.
NS:	de.	tum.de.	
DS:	de.	tum.de.	
RRSIG DS:	de.	tum.de.	
RRSIG A:			www.tum.de.
DNSKEY:		de.	tum.de.
A:			www.tum.de

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## Resolving with DNSSEC

- ❑ Resolving A **www.tum.de.**: empty resolver cache

### DNS/DNSSEC

```
Zone:      .      de.      tum.de.
NS:        de.    tum.de.
DS:        de.    tum.de.
RRSIG DS:  de.    tum.de.
RRSIG A:   .      de.      www.tum.de.
DNSKEY:    .      de.      tum.de.
A:         .      de.      www.tum.de
```

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## DNSSEC problems

- ❑ Signing and verification are mathematical complex  
need computational power
  - ❑ DNSSEC packets are larger than DNS packets
    - Larger memory footprint for servers
    - Higher network bandwidth needs
    - Larger packets ⇒
      - Fragmentation
      - Truncation
      - Fallback from UDP to TCP
- (DNSSEC requires min pkt size: 1220 Bytes (512 DNS)  
recommends pkt size: 4000 Bytes

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## DNSSEC overhead

Type	Overhead (bytes)	Comment
DNSKEY	18 + key size	RSA or ECC
DS	36	SHA-1 digest
RRSIG	46 + key size +  zone  70 +  zone	RSA ECC
NSEC	23 +  name  +  label	

- Typical key sizes in bits:
  - RSA: 1024, 1200 [Kolkman, Gieben, 2004]
  - ECC: 136, 144 [Schroepel, Eastlake, 2004]

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## Results: DNSSEC overhead

Type	Count		DNS size		DNSSEC factor		
	All	Norm	All	Norm	RSA		ECC
					All	Norm	all
Query	32.8M	5.1M	1.5G	0.3G	<b>1.1</b>	<b>1.1</b>	<b>1.1</b>
noErr	20.0M	4.2M	3.7G	0.7	<b>4.1</b>	<b>5.3</b>	<b>2.3</b>
Final	6.8M	2.5M	1.2G	0.5G	6.2	5.7	3.0
Referral	10.9M	1.3M	2.3	0.2G	2.0	2.6	1.6
Empty	2.2M	390K	.2G	44M	<b>11.7</b>	<b>10.6</b>	<b>5.2</b>
NXDomain	1.4M	500K	.2G	57M	<b>12.7</b>	<b>12.9</b>	<b>6.2</b>

- Queries: almost no overhead
- Answers: Final and Referral OK but Empty, NXDomain expensive
- Answers: RSA about twice as expensive as ECC
- Answers: All about the same as normalized

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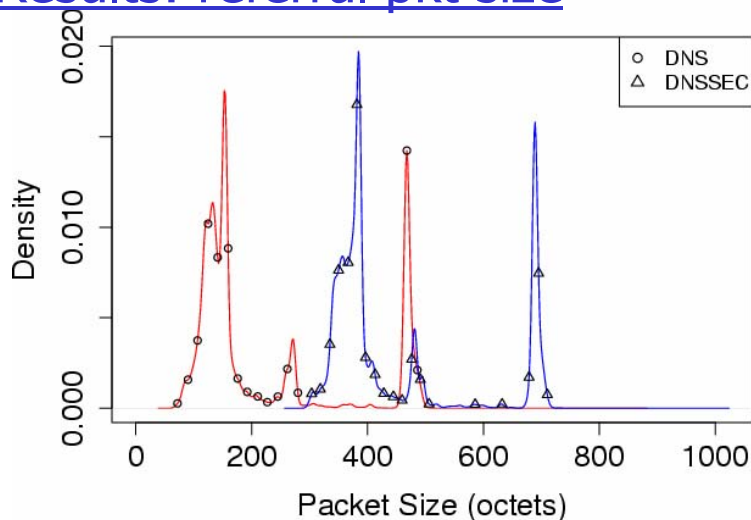
## Results: packet size

	Size ≤	NXDom	noErr	Final	Ref.	Empty
RSA	1,228	<b>.005</b>	.790	<b>.701</b>	1	<b>.633</b>
RSA	1,480	<b>.231</b>	.951	<b>.921</b>	1	<b>.996</b>
RSA	2,056	.999	.997	.991	1	.999
RSA	4,008	1.000	.999	.999	1	1.000
ECC	1,228	.998	.999	.998	1	.999
ECC	1,480	.999	.999	.999	1	.999

- ❑ Dig shows default 2056 bytes, 4008 bytes recommended
- ❑ 60% of packets not limited by maximum packet sizes (Query, FormErr, ServFail, Refused packets) but less than 1/3 of DNS size and 1/10th of DNSSEC
- ❑ noError, NXDomain problematic  
NXDomain: 77% (TR02), 72% (TR04) need fragmentation
- ❑ ECC: almost no problem...

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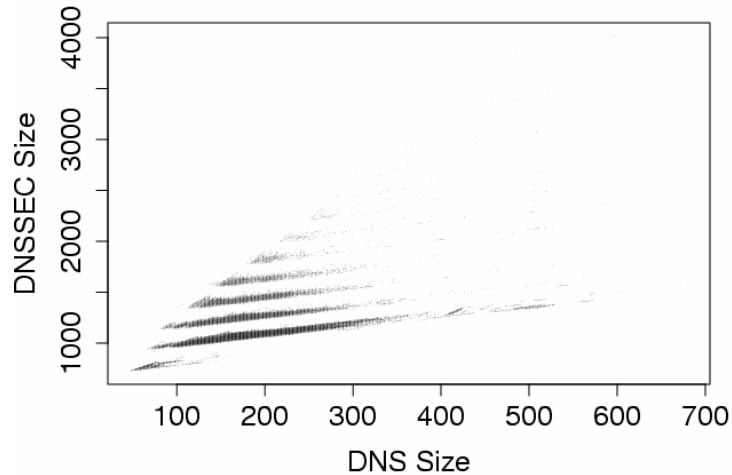
## Results: referral pkt size



- ❑ Shapes of densities similar but shifted (+1 DS RR)

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## Results: DNS vs. DNSSEC size



- Parallel lines: vertical distance 174 bytes  
same number of authoritative RRsets

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## CPU usage increase

- Authoritative name servers
  - Due to larger data volume
  - Average CPU usage (five experiments at 1 M queries)

Level	#DNS	CPU-time DNS	#DNSSEC	CPU-time DNSSEC
0	5.0K	1.0s	6.4K	2.03s
1	61.8K	11.56s	65.4K	13.85s
2	176.1K	23.80s	249.1K	39.71s

- Overhead factor 1.1 to 2
- Overhead factor per query 1.3 to 1.6

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## CPU usage increase

### ❑ Caching name servers

- Due to verification, larger data, stripping of info
- 218k queries (averaged over 5 experiments)

cached	CPU-time DNS	CPU-time DNSSEC	Delay DNS	Delay DNSSEC
No	292.8s	665.4s	1.8ms	3.9ms
Yes	37.1s	45.9s	0.2ms	0.2ms

- Without caching: factor 2.3
- Absolute additional delay small
- With caching: factor 1.25

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## Summary

- ❑ Examined cost of wide spread DNSSEC deployment
  - Network bandwidth
  - Resource consumption
- ❑ Real world analysis from large client population
- ❑ Findings:
  - Increase per packets: avg: 3.4 max: 12.7
  - ECC outperforms RSA
  - Higher memory and CPU requirements
- ❑ **No apparent show stopper!**

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