

Overlay/Underlay Interaction

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Outline

- ❑ Motivation
 - Problems with the underlying routing system
 - Source routing, overlay networks, and hybrids
- ❑ Overlay networks
 - Pros: flexibility, limited overhead, & value-added
 - Cons: data-path overhead, probes, & feedback
- ❑ Negative interactions
 - With other overlays: the price of anarchy
 - With the underlay: influence on traffic engineering
 - With itself: bi-stability and trunk reservation
- ❑ Future directions

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What's Wrong With Internet Routing?

- ❑ Restrictive path-selection model
 - Destination-based packet forwarding
 - Single best BGP path per prefix
 - BGP routing constrained by policies
 - Ignoring congestion and delay
 - Ignoring application requirements
- ❑ Unappealing protocol dynamics
 - Persistent oscillation (due to policy conflicts)
 - Slow convergence (due to path exploration)
 - Lost reachability (due to route-flap damping)

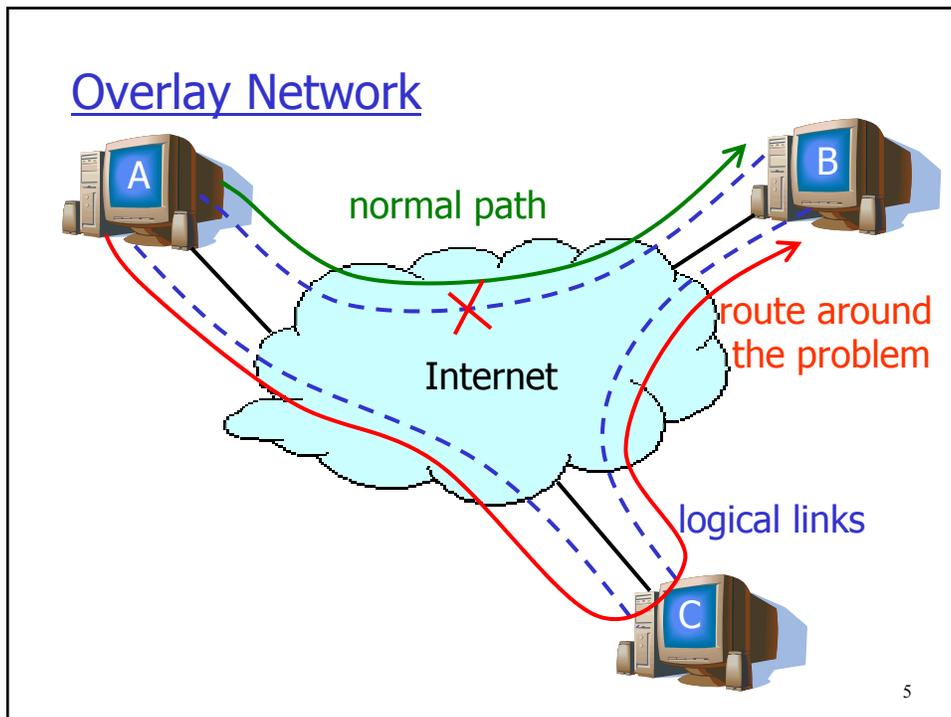
Need for routing to scale to millions of routers

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Putting More Power in End Hosts

- ❑ Source routing (e.g., Nimrod)
 - End host selects the end-to-end path
 - Routers simply forward packets on the path
 - Requires the routers to agree to participate
- ❑ Overlay networks (e.g., RON)
 - Conventional computers act as logical routers
 - Real routers deliver packets to intermediate hosts
 - No need for cooperation from the real routers
- ❑ Hybrid schemes
 - Source routing at the AS level
 - Source routing in the overlay network

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- ### Advantage: Flexible Routing
- ❑ Paths that violate BGP routing policy
 - E.g., A to C goes through AT&T and Sprint
 - ... and C to B goes through UUNET
 - BGP would not allow AT&T-Sprint-UUNET path
 - ❑ Quick adaptation to network problems
 - Fast detection of congestion and outages
 - ... by probing as aggressively as necessary
 - ❑ Selecting paths based on different metrics
 - E.g., overlay selects paths based on latency
 - ... whereas the underlay might try to balance load
- The number 6 is in the bottom right corner.

Advantage: Fewer Worries About Scalability

- ❑ Small number of nodes
 - Just enough nodes to have diverse paths
 - A few friends who want better service
 - Virtual Private Network of several corporate sites
- ❑ Balancing the trade-offs
 - High probe frequency for maximum adaptivity
 - Low probe frequency for minimum overhead
- ❑ Simple routing protocol
 - Link-state protocol to learn probing results
 - Selecting a good intermediate hop when needed

Deploy multiple small overlay networks, if necessary₇

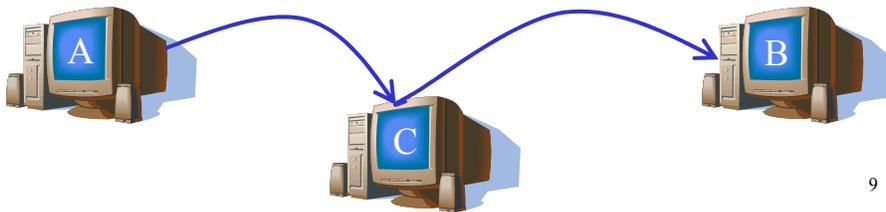
Advantage: Customizing Packet Delivery

- ❑ Recovering from packet loss
 - Packet retransmission
 - Forward error correction
- ❑ Quality-of-service differentiation
 - Classify packets based on header bits
 - Schedule packet transmissions based on result
- ❑ Incremental deployment of new features
 - Multicast communication (e.g., MBone)
 - IPv6 (e.g., 6Bone)
 - Encryption of packet contents

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Disadvantage: Traversing Intermediate Nodes

- ❑ Processing delay
 - Packets going through multiple software nodes
- ❑ Network performance
 - Propagation delay on circuitous path
 - Network congestion from extra load
- ❑ Financial cost
 - Bill for traffic going in/out of intermediate node



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Disadvantage: Limitations of Active Probes

- ❑ Bandwidth overhead
 - Probe traffic between two nodes
 - Propagating probe results to other nodes
- ❑ Limited accuracy of end-to-end probes
 - Available bandwidth of logical link?
 - Losses due to congestion vs. failure?
 - Problem on forward vs. reverse path?
- ❑ Limited visibility
 - Logical links may share underlay routers/links
 - May be hard to detect the dependencies

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Disadvantage: Feedback Effects

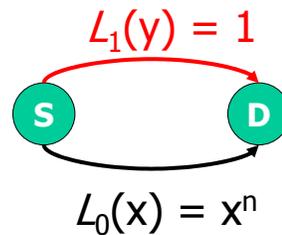
- Background traffic
 - Overlay traffic consumes extra resources
 - ... at the expense of regular background traffic
 - But, the overlay traffic *does* get out of the way!
- Other overlays
 - Potential competition between multiple overlays
 - E.g., one overlay picks a (longer) alternate path
 - ... and extra load causes another overlay to adapt
- Underlying network
 - Overlay network changes the traffic matrix
 - ... forcing operators to adapt the underlay routing

Are these effects significant? Any positive effects?

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Price of Anarchy (Roughgarden&Tardos)

- Worst-case example
 - Two paths from s to d
 - Total of one unit of load
 - Latency as function of load
- Selfish source routing
 - All traffic through bottom link
 - Mean latency of 1
- Latency-optimal routing
 - Minimize mean latency
 - Set $x = [1/(n+1)]^{1/n}$
 - Mean latency goes to 0



Total load: $x + y = 1$
Mean latency: $x L_0(x) + y L_1(y)$

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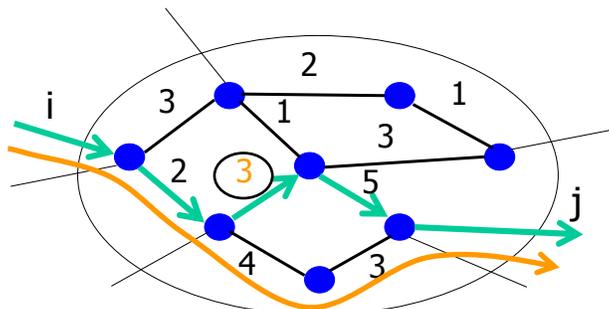
Internet-like Environments (Qiu et al.)

- Realistic networks
 - Backbone network topologies
 - Link delay (propagation and queuing delay)
 - Routing set to minimize network congestion
- Realistic overlays
 - Small number of overlay nodes (limited flexibility)
 - Overlay paths chosen to minimize latency
- Practice doesn't match the worst-case theory
 - Some tension between the two different metrics
 - But, not anywhere near as bad as the worst case

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Interaction With Traffic Eng. (Qiu et al)

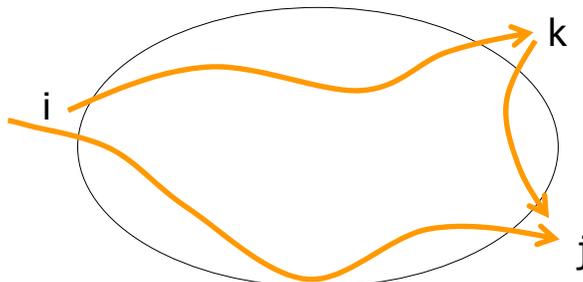
- Underlay network: traffic engineering
 - Inputs: **traffic matrix** M_{ij} and physical topology
 - Objective: minimize overall network congestion
 - Output: **routing** R_{ij} : fraction of (i,j) traffic on link l
 - Propagation and queuing **delay on virtual links**



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Interaction With Traffic Eng. (Qiu et al)

- Overlay network: selects intermediate nodes
 - Inputs: measured **delay for each virtual link**
 - Objective: minimizing end-to-end latency
 - Output: choice of intermediate nodes for traffic
 - **Traffic matrix** on the underlay network



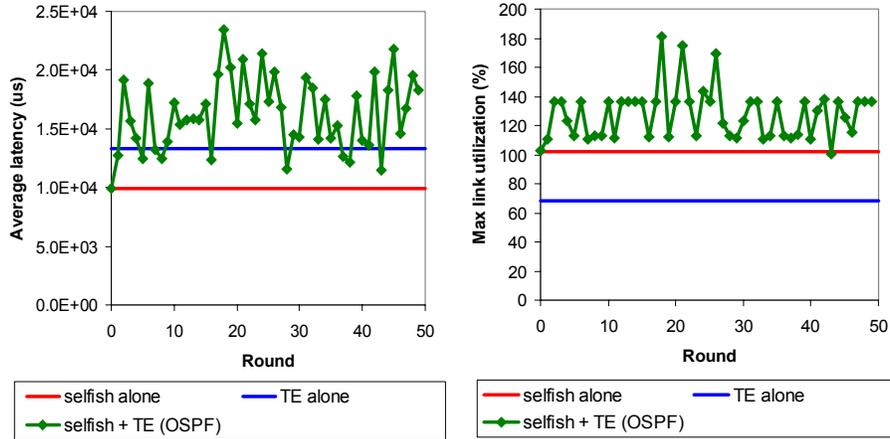
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Interaction With Traffic Eng. (Qiu et al)

- Underlay network: traffic engineering
 - Inputs: **traffic matrix** and physical topology
 - Objective: minimize overall network congestion
 - Output: selection of paths in underlay network
 - Propagation and queuing **delay on virtual links**
- Overlay network: selecting intermediate nodes
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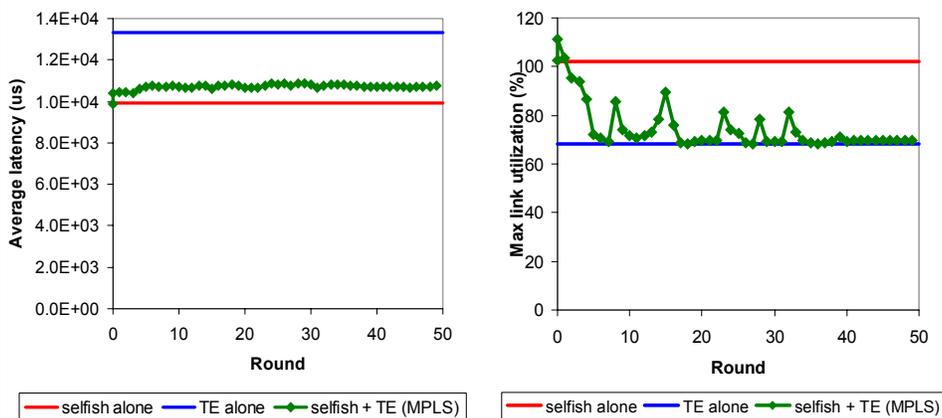
Interaction TE vs. OSPF Weight Tweaks



OSPF optimizer interacts poorly with selfish overlays because it only has very coarse-grained control.

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Interaction TE vs. Multi-commodity Flow

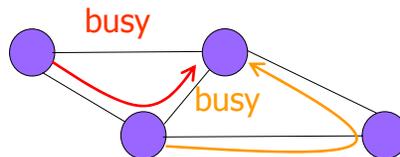


Multi-commodity flow optimizer interacts with selfish overlays much more effectively.

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History: Bistability in Single Overlay

- Phone network is an overlay
 - Logical link between each pair of switches
 - Phone call put on one-hop path, when possible
 - ... and two-hop alternate path otherwise
- Problem: inefficient path assignment
 - Two-hop path for one phone call
 - ... stops another call from using direct path
 - ... forcing the use of a two-hop alternate path



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Preventing Inefficient Routes: Trunk Reservation

- Two stable states for the system
 - Mostly one-hop calls with low blocking rate
 - Mostly two-hop calls with high blocking rate
- Making the system stable
 - Reserve a portion of each link for direct calls
 - When link load exceeds threshold...
 - ... disallow two-hop paths from using the link
 - Rejects some two-hop calls
 - ... to keep some spare capacity for future one-hop calls
- Stability through trunk reservation
 - Single efficient, stable state with right threshold

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Should ISPs Fear Overlays, or Favor Them?

- ❑ Billing
 - Con: overlays commoditize the network providers
 - Pro: overlay traffic adds traffic subject to billing
- ❑ Engineering
 - Con: traffic matrix becomes less predictable
 - Pro: TE less important because overlays can adapt
- ❑ Value-added services
 - Con: overlays become the place for new services
 - Pro: ISPs can provide overlay nodes in the core

Do the pros outweigh the cons? Beat 'em, or join 'em?

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What Could ISPs Do to Help Overlays?

- ❑ Better visibility
 - Measurements from the underlying network
- ❑ Better control
 - Influence over underlay path selection
- ❑ Lower cost
 - Avoiding need to traverse intermediate hosts
- ❑ Better joint optimization
 - Underlay adaptation accountings for the overlays

- ❑ Idea: ISP could provide an oracle to the P2P

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ISP Support: More Visibility Via Data

- ❑ Underlying topology
 - Dependencies: between paths for virtual links
 - Resource limits: capacity of the underlying links
- ❑ Routing-protocol update streams
 - Fast adaptation: BGP may provide early warning
 - Better adaptation: identify location of a problem
- ❑ Performance measurement
 - Statistics: per-link delay, loss, and throughput
 - Efficiency: consolidating probes for many overlays

Are there enough incentives to share the data?

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ISP Support: Direct Control

- ❑ Influence over paths for virtual links
 - Guarantees of link/node disjoint paths?
- ❑ Control over packet forwarding
 - Overlays allowed to install forwarding entries?
- ❑ Warnings about planned changes
 - Prior notification about planned maintenance
 - Rerouting of virtual link before the outage

Are there effective ways to share control?

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ISP Support: Lower Cost

- ❑ Measurement service to overlays
 - Reduce overhead and cost of active probing
- ❑ Deflection points inside the network
 - Avoiding packets traversing intermediate hosts
 - Encapsulation to deflect packet through landmark
- ❑ Hosting of servers for running overlays
 - Enhanced form of a hosting service
 - Servers running directly in the ISP's PoPs

Reasonable business models and technical solutions?

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ISP Support: Joint Optimization

- ❑ Changes in traffic-engineering practices
 - Overlay-friendly traffic engineering?
 - Ways to ensure stability and optimality?
 - Or, just let the overlays manage the traffic?
- ❑ Cross-domain cooperation
 - Management of paths and traffic across ASes
 - Virtual links that span two (or more) ISPs
 - Ways to ensure stability and optimality?
 - Service-level agreements spanning ASes?

Reasonable optimization or game-theory models?

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Discussion

- ❑ Are overlay networks a good idea?
 - Just a hack to avoid changing the underlay?
 - What if we could “fix” the underlying network?
 - Would we still have a need for overlay networks?
- ❑ Should we have overlay-friendly underlays?
 - Or underlay-friendly overlays, or both?
 - Visibility, control, economics, efficiency, ...
 - Or, are the two systems inherently at odds?
- ❑ What about interactions between overlays?
 - Cooperate to reduce measurement cost and prevent suboptimality and instability?
 - Compete because that’s the way life works?

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Conclusions

- ❑ Overlays
 - Enables innovation in routing and forwarding
 - ... without changing the underlying network
- ❑ Interaction effects
 - With background traffic
 - With other overlays
 - With traffic engineering
- ❑ Ongoing work
 - Right interplay between underlay and overlay?

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