

## Network layer: Overview

- ❑ Network layer functions
- ❑ IP
- ❑ Routing and forwarding
- ❑ NAT
- ❑ ARP
- ❑ Routing

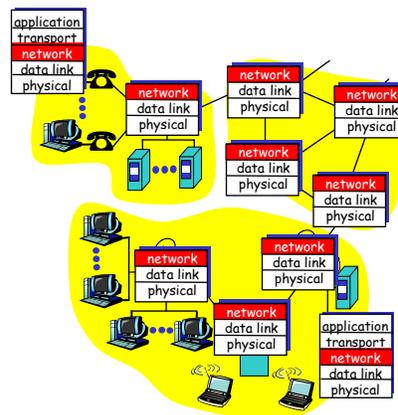
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## Network layer functions

- ❑ Transport packet from sending to receiving hosts
- ❑ Network layer protocols in *every* host, router

### Three important functions:

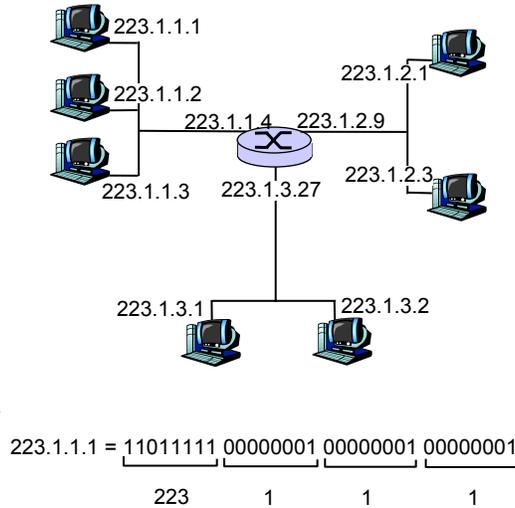
- ❑ *Path determination*: route taken by packets from source to dest. *Routing algorithms*
- ❑ *Switching*: move packets from router's input to appropriate router output



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## IP addressing

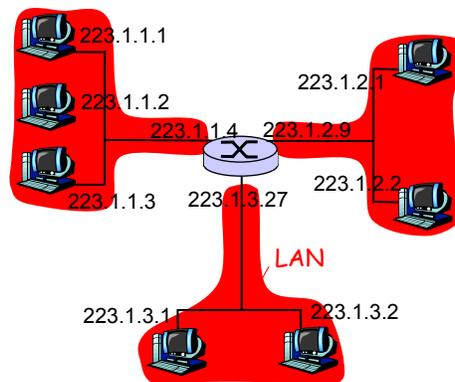
- **IP address:** 32-bit identifier for host, router *interface*
- **Interface:** connection between host, router and physical link
  - router's typically have multiple interfaces
  - host may have multiple interfaces
  - IP addresses associated with interface, not host, router



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## IP addressing (2)

- **IP address:**
  - Network part (high order bits)
  - Host part (low order bits)
- **What's a network ?**  
(from IP address perspective)
  - Device interfaces with same network part of IP address
  - Can physically reach each other without intervening router



Network consisting of 3 IP networks  
(for IP addresses starting with 223,  
first 24 bits are network address)

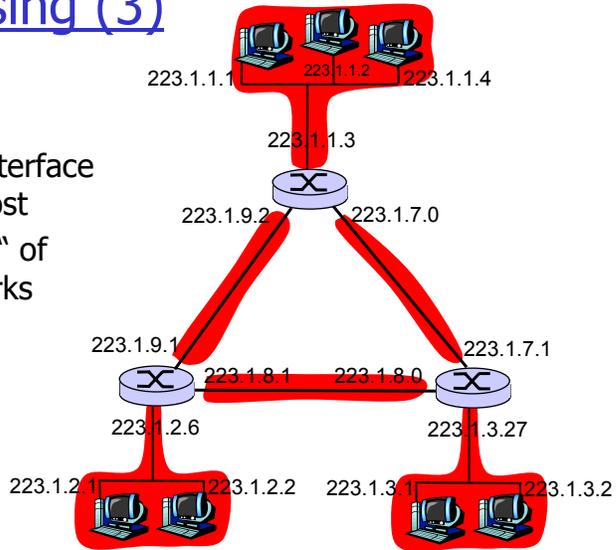
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## IP addressing (3)

How to find the networks?

- Detach each interface from router, host
- Create "islands" of isolated networks

Interconnected system consisting of six networks



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## IP networks: Subnets

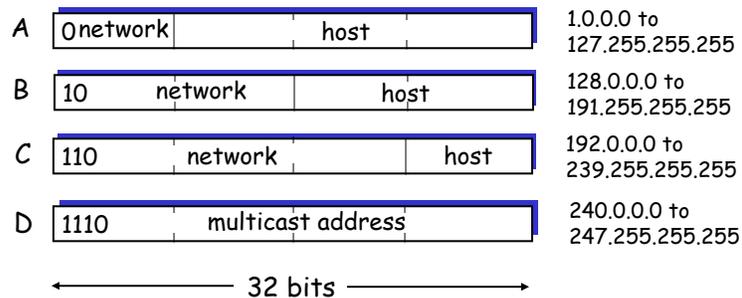
- Sub divide address space
  - network part
  - host address
- Address format: **a.b.c.d/x**, where x is # bits in subnet portion of address



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## Fixed subnetting (classful)

class



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## Address management

- Problem: we are running out of networks
- Solution (a):  
**subnetting:** e.g., Class B Host field (16bits) is subdivided into <subnet;host> fields
- Solution (b):  
**CIDR** (Classless Inter Domain Routing)

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## CIDR

- ❑ Classless InterDomain Routing
- ❑ Motivation
  - Class A is too large, Class C is too small
  - Everyone had a Class B address!!!
- ❑ Solution:
  - Sites are given contiguous blocks of class-C addresses (256 addresses each) and a mask or parts of former class A/B networks.

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## CIDR (2.)

### CIDR: Classless InterDomain Routing

- ❑ Subnet portion of address of arbitrary length
- ❑ Address format: **a.b.c.d/x**, where x is # bits in subnet portion of address



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## IP addresses: how to get one?

**Q:** How does *host* get IP address?

- ❑ Hard-coded by system admin in a file
  - Wintel: control-panel->network->configuration->tcp/ip->properties
  - UNIX: /etc/rc.config
- ❑ **DHCP: Dynamic Host Configuration Protocol:** dynamically get address from as server
  - "Plug-and-play"
- ❑ IP / Subnets allocated by provider (RIPE/ARIN/...)

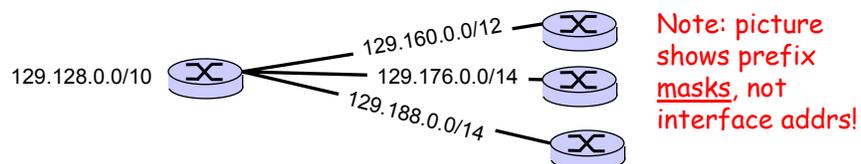
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## Hierarchical address structure

❑ Recall: CIDR

128.119.48.12/18 ← 10000000 01110111 00110000 00001100

- High order bits form the **prefix**
- Once inside the network, can **subnet**: divide remaining bits
- Subnet example:



❑ **Forwarding decision: longest prefix match**

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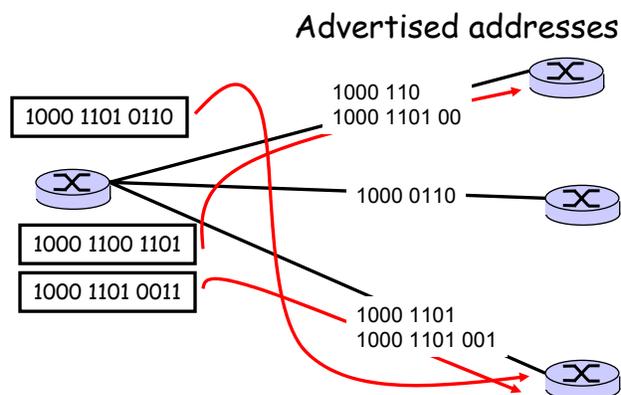
## Forwarding vs. routing

- **Forwarding**: the process of moving packets from input to output
    - The forwarding table
    - Information in the packet
  
  - **Routing**: process by which the forwarding table is built and maintained
    - One or more routing protocols
    - Procedures (algorithms) to convert routing info to forwarding table.
- (More later....)

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## Forwarding with CIDR

- Packet should be sent toward the interface with the **longest matching prefix**



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## CIDR (3.)

- ❑ Repeated “aggregation” within same provider leads to shorter and shorter prefixes
- ❑ CIDR helps also routing table size and processing: Gateways keep only prefixes and find “longest prefix” match
- ❑ Class-C is also partitioned by geography e.g., Europe got 194.0.0.0 to 195.255.255.255

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## Lookup: Longest prefix match

- ❑ Forwarding table:  
*<Network>/<mask> <next-hop>*
- ❑ IP Packets: destination IP address
  - Find next-hop via longest prefix match
- ❑ Example:

Forwarding table		Packets
134.96.252.0/24	A	134.96.252.200
134.96.0.0/16	C	134.96.254.2
134.96.240.0/20	B	134.96.239.200
134.96.252.192/28	B	134.97.239.200
134.96.252.128/28	A	134.96.252.191

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## IP addressing: The last word...

**Q:** How does an ISP get block of addresses?

**A:** **ICANN:** Internet Corporation for Assigned Names and Numbers

- allocates addresses
- manages DNS
- assigns domain names, resolves disputes

**Q:** What do I do if I don't have a public address?

**A:** **Private IP addresses** (RFC1918)

- 10/8
- 172.16/12
- 192.168/16
- Private use only – not routable in the Internet

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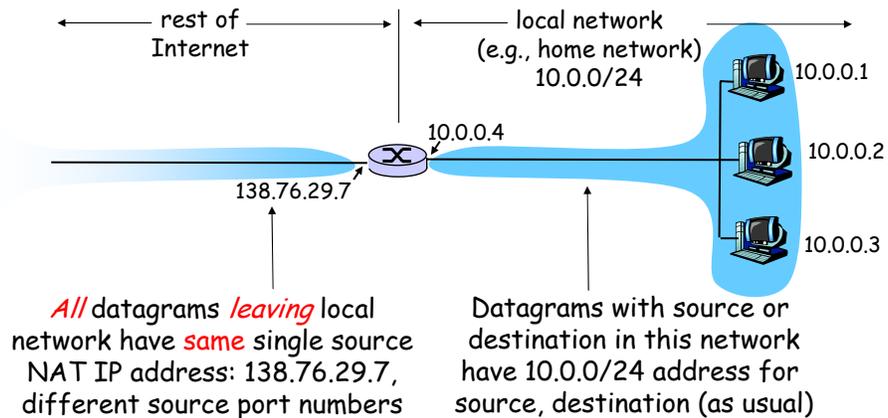
## NAT: Network Address Translation

**Motivation:** local network uses just one IP address as far as outside world is concerned:

- Just one IP address for all devices
- Not needed range of addresses from ISP

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## NAT: Network Address Translation



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## NAT: Network Address Translation

**Motivation:** Local network uses just one IP address as far as outside world is concerned:

- Range of addresses not needed from ISP: just one IP address for all devices
- Can change addresses of devices in local network without notifying outside world
- Can change ISP without changing addresses of devices in local network
- Devices inside local net not explicitly addressable, visible by outside world (a security plus).

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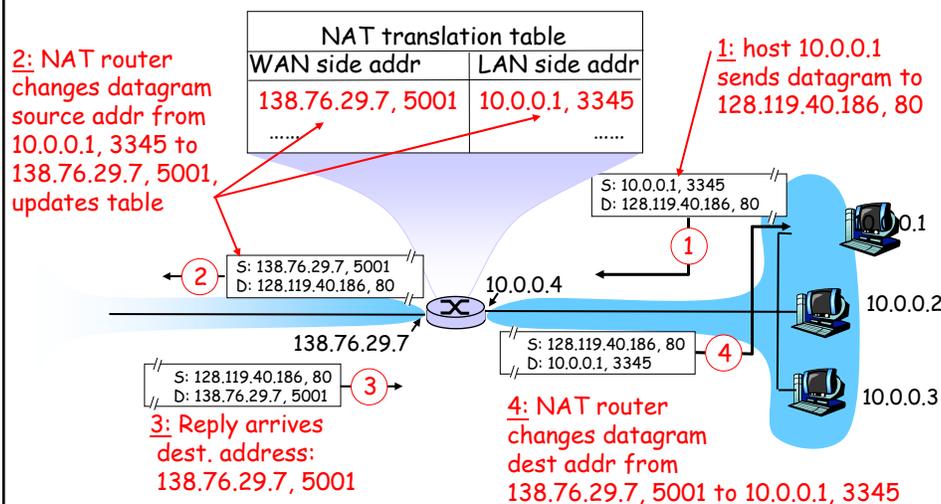
## NAT: Network Address Translation

**Implementation:** NAT router must:

- *Outgoing datagrams: replace* (source IP address, port #) of every outgoing datagram to (NAT IP address, new port #)
  - . . . remote clients/servers will respond using (NAT IP address, new port #) as destination addr.
- *Remember (in NAT translation table)* every (source IP address, port #) to (NAT IP address, new port #) translation pair
- *Incoming datagrams: replace* (NAT IP address, new port #) in dest fields of every incoming datagram with corresponding (source IP address, port #) stored in NAT table

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## NAT: Network Address Translation



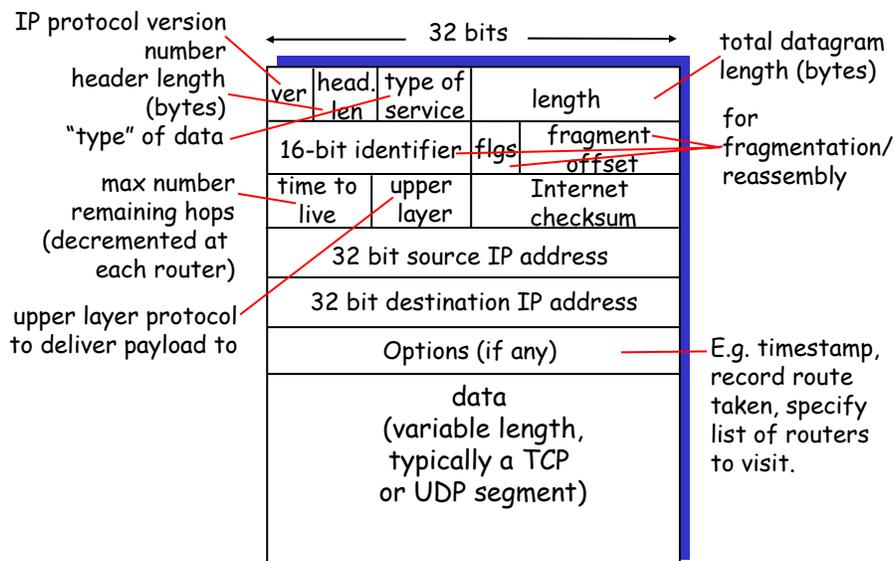
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## NAT: Network Address Translation

- 16-bit port-number field:
  - 60,000 simultaneous connections with a single LAN-side address!
- NAT is controversial:
  - Routers should only process up to layer 3
  - Violates end-to-end argument
    - NAT possibility must be taken into account by app designers, eg, P2P applications
  - Address shortage should instead be solved by IPv6

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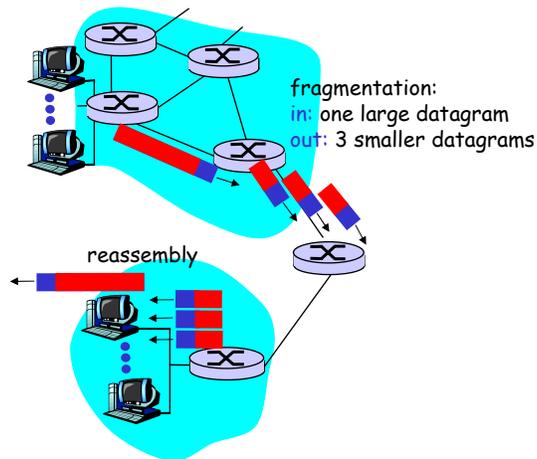
## IPv4 datagram format



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## IP fragmentation and reassembly

- Network links have MTU (max.transfer size) - largest possible link-level frame.
  - Different link types, different MTUs
- Large IP datagram divided ("fragmented") within net
  - One datagram becomes several datagrams
  - "Reassembled" only at final destination
  - IP header bits used to identify, order related fragments



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## IP fragmentation and reassembly (2.)

### Example

- 4000 byte datagram
- MTU = 1500 bytes

length	ID	fragflag	offset
=4000	=x	=0	=0

One large datagram becomes several smaller datagrams

1480 bytes in data field

offset = 1480/8

length	ID	fragflag	offset
=1500	=x	=1	=0
=1500	=x	=1	=185
=1040	=x	=0	=370

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## ICMP: Internet Control Message Protocol

- ❑ Used by hosts, routers, gateways to communication network-level information
    - Error reporting: unreachable host, network, port, protocol
    - Echo request/reply (used by ping)
  - ❑ Network-layer "above" IP:
    - ICMP msgs carried in IP datagrams
  - ❑ ICMP message: type, code plus first 8 bytes of IP datagram causing error
- | Type | Code | description                                   |
|------|------|---|
| 0    | 0    | echo reply (ping)                             |
| 3    | 0    | dest. network unreachable                     |
| 3    | 1    | dest host unreachable                         |
| 3    | 2    | dest protocol unreachable                     |
| 3    | 3    | dest port unreachable                         |
| 3    | 6    | dest network unknown                          |
| 3    | 7    | dest host unknown                             |
| 4    | 0    | source quench (congestion control - not used) |
| 8    | 0    | echo request (ping)                           |
| 9    | 0    | route advertisement                           |
| 10   | 0    | router discovery                              |
| 11   | 0    | TTL expired                                   |
| 12   | 0    | bad IP header                                 |

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## Address Resolution Protocol (ARP)

- ❑ Interface between Link layer - Network Layer
- ❑ Allows hosts to query who owns an IP address on the same LAN
- ❑ Owner responds with hardware address
- ❑ Allows changes to link layer to be independent of IP addressing

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## Network layer: status

- Network layer functions
- IP
- Routing and forwarding
- NAT
- ARP
- Routing