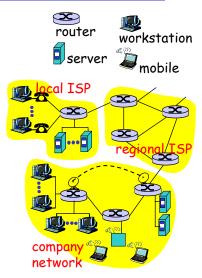
# Network Protocols and Architectures

#### Introduction

1

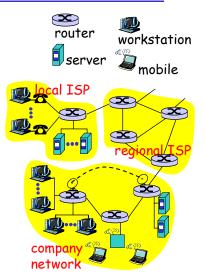
#### What's the Internet: "nuts and bolts" view

- Millions of connected computing devices: hosts, end-systems
  - PC's workstations, servers
  - PDA's, phones, toasters running network apps
- Communication links
  - Fiber, copper, radio, satellite
- Routers: forward packets (chunks) of data through network



#### What's the Internet: "nuts and bolts" view

- Protocols: control sending, receiving of messages
  - o E.g., TCP, IP, HTTP, FTP, PPP
- Internet: "network of networks"
  - Loosely hierarchical
  - Public Internet versus private intranet
- Internet standards
  - RFC: Request for comments
  - IETF: Internet Engineering Task Force

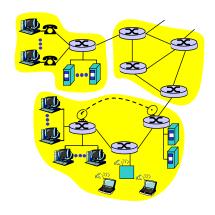


3

## What's the Internet: A service view

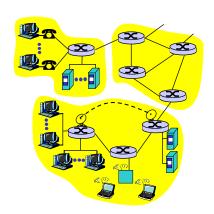
- Communication *infrastructure* enables distributed applications:
  - WWW, email, games, ecommerce, database, voting,
  - More?
- Communication services provided:
  - Connectionless
  - Connection-oriented
- cyberspace [Gibson]:

"a consensual hallucination experienced daily by billions of operators, in every nation, ...."



## A closer look at network structure

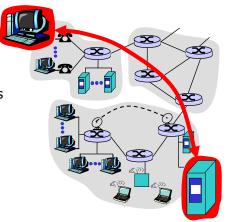
- Network edge: applications and hosts
- □ Network core:
  - Routers
  - Network of networks
- Access networks, physical media: Communication links



5

## The network edge

- End systems (hosts):
  - Run application programs
  - o E.g., WWW, e-mail
  - At "edge of network"
- Client/server model
  - Client host requests, receives services from server
  - E.g., WWW client (browser)/ server; e-mail client/server
- □ Peer-peer model:
  - Host interaction symmetric
  - E.g., teleconferencing



## Network edge: connection-oriented service

- *Goal:* data transfer between end sys.
- Handshaking: setup (prepare for) data transfer ahead of time
  - Hello, hello back human protocol
  - Set up "state" in two communicating hosts
- TCP Transmission Control Protocol
  - Internet's connectionoriented service

#### TCP service [RFC 793]

- Reliable, in-order bytestream data transfer
  - Loss: acknowledgements and retransmissions
- □ Flow control:
  - Sender won't overwhelm receiver
- Congestion control:
  - Senders "slow down sending rate" when network congested

7

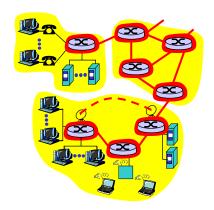
## Network edge: connectionless service

Goal: Data transfer between end systems

- Same as before!
- □ UDP User Datagram Protocol [RFC 768]: Internet's connectionless service
  - Unreliable data transfer
  - No flow control
  - No congestion control

## The network core

- Mesh of interconnected routers
- The fundamental question: How is data transferred through net?
  - Circuit switching:
    Dedicated circuit per call: telephone net
  - Packet switching: Data sent through net in discrete "chunks"

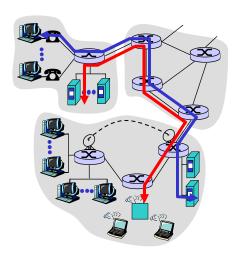


9

## Network core: Circuit switching

## End-end resources reserved for "call"

- Link bandwidth, switch capacity
- Dedicated resources: no sharing
- Circuit-like (guaranteed) performance
- Call setup required



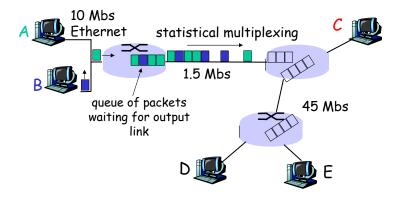
## Network core: Packet switching

#### Each end-end data stream divided into packets

- ☐ Users' A, B packets *share* network resources
- □ Each packet uses full link bandwidth
- □ Resources used *as needed*

11

## Network core: Packet switching



Packet-switching versus circuit switching: Human restaurant analogy

## Network core: Packet switching

#### Resource contention:

- ☐ Aggregate resource demand can exceed amount available
- Congestion: packets queue, wait for link use
- ☐ Store and forward: packets move one hop at a time
  - Transmit over link
  - Wait turn at next link

13

## Packet switching vs. circuit switching

Is packet switching a "slam dunk winner?"

- Great for bursty data
  - Resource sharing
  - No call setup
- Excessive congestion: packet delay and loss
  - Protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
  - Bandwidth guarantees needed for audio/video apps still an unsolved problem

## Packet-switched networks: Routing

- Goal: Move packets among routers from source to destination
  - We'll study several path selection algorithms
- Datagram network:
  - Destination address determines next hop
  - Routes may change during session
  - Analogy: driving, asking directions
- Virtual circuit network:
  - Each packet carries tag (virtual circuit ID), tag determines next hop
  - Fixed path determined at call setup time, remains fixed through call
  - Routers maintain per-call state

15

## Protocol "layers"

#### Networks are complex!

- Many "pieces":
  - Hosts
  - Routers
  - Links of various media
  - Applications
  - Protocols
  - Hardware, software

#### **Question:**

Is there any hope of organizing structure of network?

Or at least in our discussion of networks?

## Why layering?

#### Dealing with complex systems:

- Explicit structure allows identification, relationship of complex system's pieces
  - Layered reference model for discussion
- Modularization eases maintenance, updating of system
  - Change of implementation of layer's service transparent to rest of system
  - E.g., change in gate procedure does not affect rest of system
- Layering considered harmful?

17

## <u>Internet protocol stack</u>

- Application: supporting network applications
- ☐ Transport: host-host data transfer
- Network: uniform format of packets, routing of datagrams from source to destination
- Link: data transfer between neighboring network elements
- Physical: bits "on the wire"

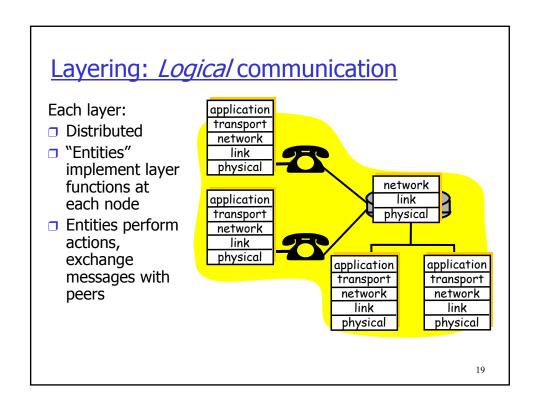
application

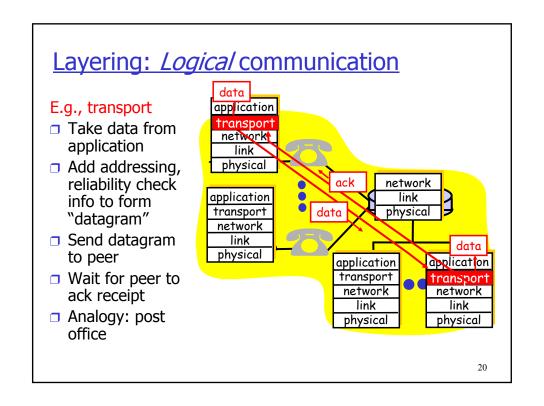
transport

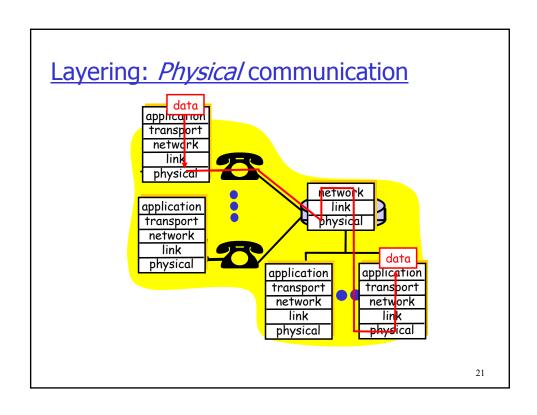
network

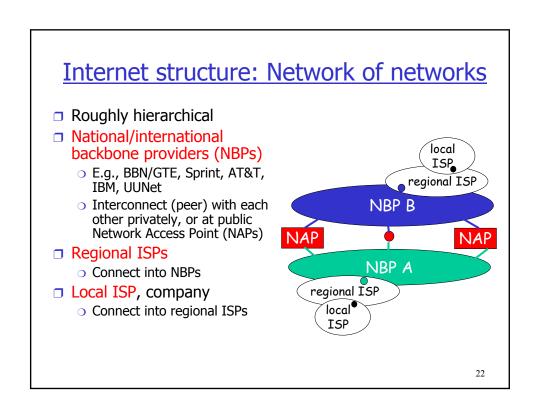
link

physical









## Principles of the Internet

- ☐ Edge vs. core (end-systems vs. routers)
  - Dumb network
  - Intelligence at the end-systems
- Different communication paradigms
  - Connection oriented vs. connection less
  - Packet vs. circuit switching
- Layered System
- □ Network of collaborating networks