

Network layer: Overview

- ❑ Network layer functions
- ❑ IP
- ❑ Routing and forwarding
- ❑ NAT
- ❑ ARP
- ❑ IPv6
- ❑ Routing

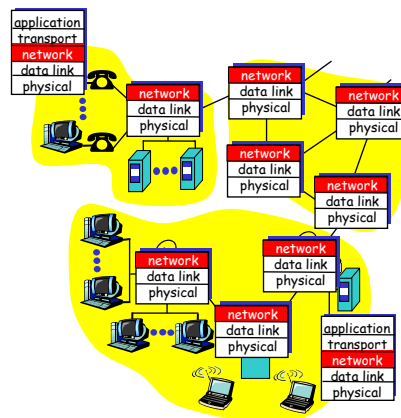
1

Network Layer Functions

- ❑ Transport packet from sending to receiving hosts
- ❑ Network layer protocols in *every* host, router

Three important functions:

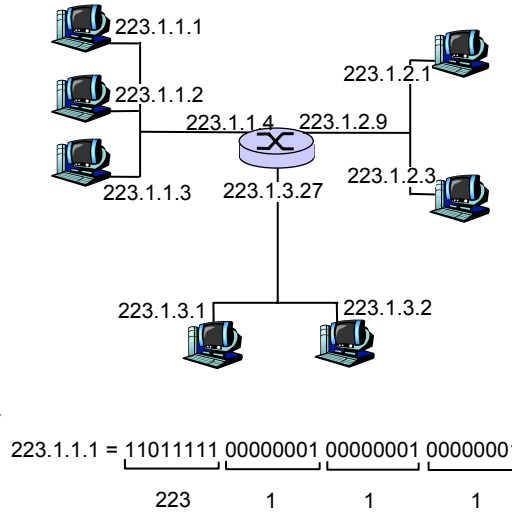
- ❑ *Path determination*: route taken by packets from source to dest. *Routing algorithms*
- ❑ *Switching*: move packets from router's input to appropriate router output



2

IP addressing

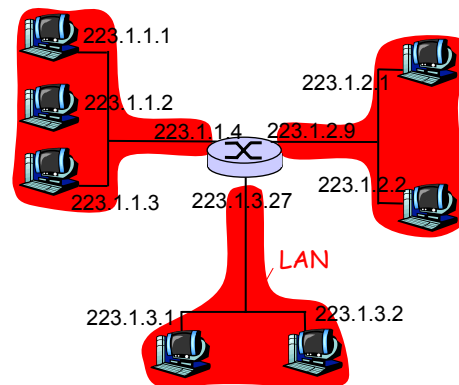
- IP address: 32-bit identifier for host, router *interface*
- *Interface*: connection between host, router and physical link
 - routers typically have multiple interfaces
 - host may have multiple interfaces
 - IP addresses associated with interface, not host, router



3

IP addressing (2)

- IP address:
 - Network part (high order bits)
 - Host part (low order bits)
- *What's a network?*
(from IP address perspective)
 - Device interfaces with same network part of IP address
 - Can physically reach each other without intervening router



Network consisting of 3 IP networks
(for IP addresses starting with 223,
first 24 bits are network address)

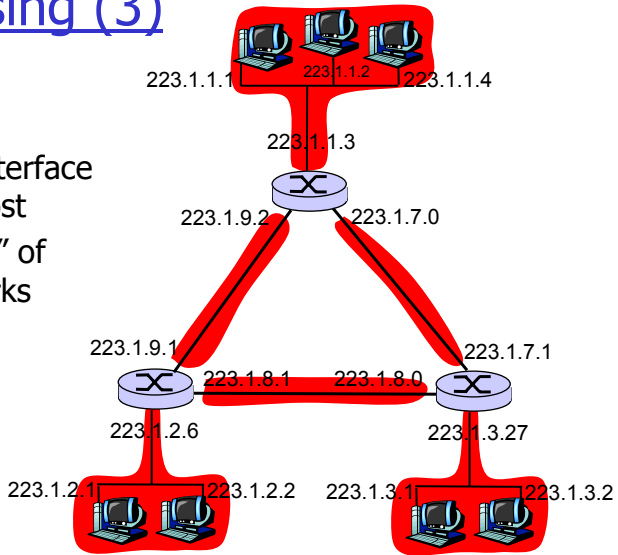
4

IP addressing (3)

How to find the networks?

- ❑ Detach each interface from router, host
- ❑ Create "islands" of isolated networks

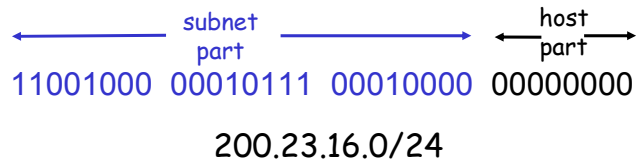
Interconnected system consisting of six networks



5

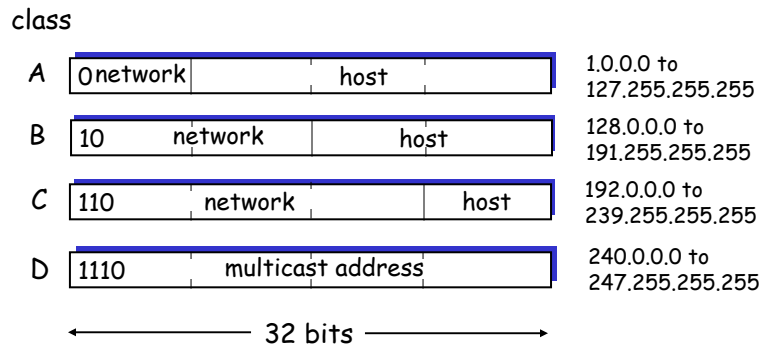
IP networks: Subnets

- ❑ Sub divide address space
 - network part
 - host address
- ❑ Address format: **a.b.c.d/x**, where x is # bits in subnet portion of address



6

Fixed subnetting (classful)



7

Address management

- Problem: We are running out of networks
- Solution (a):
Subnetting: e.g., Class B Host field (16 bits) is subdivided into <subnet;host> fields
- Solution (b):
CIDR (Classless Inter Domain Routing)

8

CIDR

CIDR: Classless InterDomain Routing

□ Motivation

- Class A is too large, Class C is too small
- Everyone had a Class B address!!!

□ Solution:

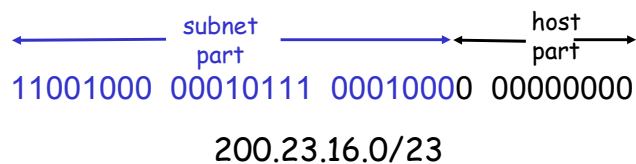
- Sites are given contiguous blocks of class-C addresses (256 addresses each) and a mask or parts of former class A/B networks.

9

CIDR (2.)

CIDR: Classless InterDomain Routing>

- Subnet portion of address of arbitrary length
- Address format: **a.b.c.d/x**, where x is # bits in subnet portion of address



10

IP addresses: How to get one?

Q: How does *host* get IP address?

- Hard-coded by system admin in a file
 - Wintel: Control Panel → Network → Configuration → TCP/IP → Properties
 - UNIX: /etc/rc.config
- **DHCP: Dynamic Host Configuration Protocol:** dynamically get address from as server
 - "Plug-and-play"
- IP / Subnets allocated by provider (RIPE/ARIN/...)

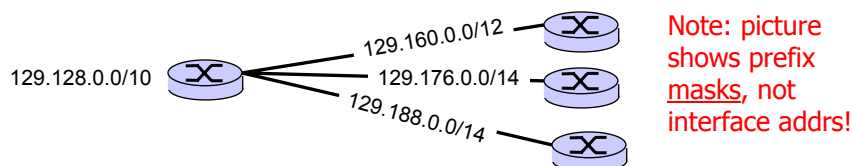
11

Hierarchical address structure

□ Recall: CIDR

128.119.48.12/18 = 10000000 01110111 00110000 00001100

- High order bits form the **prefix**
- Once inside the network, can **subnet**: divide remaining bits
- Subnet example:



□ **Forwarding decision: Longest prefix match**

12

Forwarding vs. routing

- **Forwarding:** the process of moving packets from input to output
 - The forwarding table
 - Information in the packet

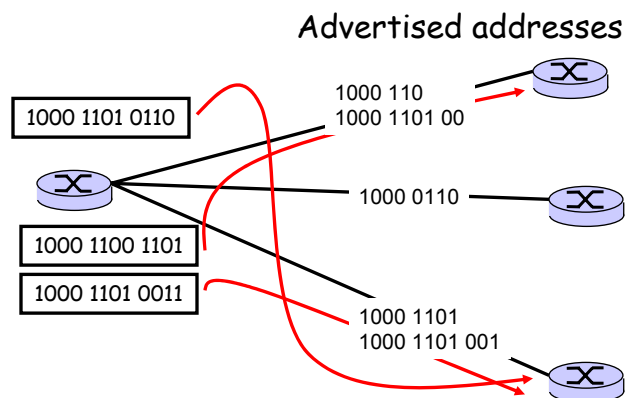
- **Routing:** process by which the forwarding table is built and maintained
 - One or more routing protocols
 - Procedures (algorithms) to convert routing info to forwarding table.

(More later ...)

13

Forwarding with CIDR

- Packet should be sent toward the interface with the **longest matching prefix**



14

CIDR (3.)

- ❑ Repeated “aggregation” within same provider leads to shorter and shorter prefixes
- ❑ CIDR helps also routing table size and processing: Gateways keep only prefixes and find “longest prefix” match
- ❑ Class-C networks are also partitioned by geography e.g., Europe got 194.0.0.0 to 195.255.255.255

15

Lookup: Longest prefix match

- ❑ Forwarding table:
<Network>/<mask> <next-hop>
- ❑ IP Packets: destination IP address
 - Find next-hop via longest prefix match
- ❑ Example:

| Forwarding table | | Packets |
|-------------------|---|----------------|
| 134.96.252.0/24 | A | 134.96.252.200 |
| 134.96.0.0/16 | C | 134.96.254.2 |
| 134.96.240.0/20 | B | 134.96.239.200 |
| 134.96.252.192/28 | B | 134.97.239.200 |
| 134.96.252.128/28 | A | 134.96.252.191 |

16

IP addressing: The last word ...

Q: How does an ISP get block of addresses?

A: **ICANN:** Internet Corporation for Assigned Names and Numbers

- allocates addresses
- manages DNS
- assigns domain names, resolves disputes

Q: What do I do if I don't have a public address?

A: **Private IP addresses** (RFC 1918)

- 10/8
- 172.16/12
- 192.168/16

Private use only – not routable in the Internet

17

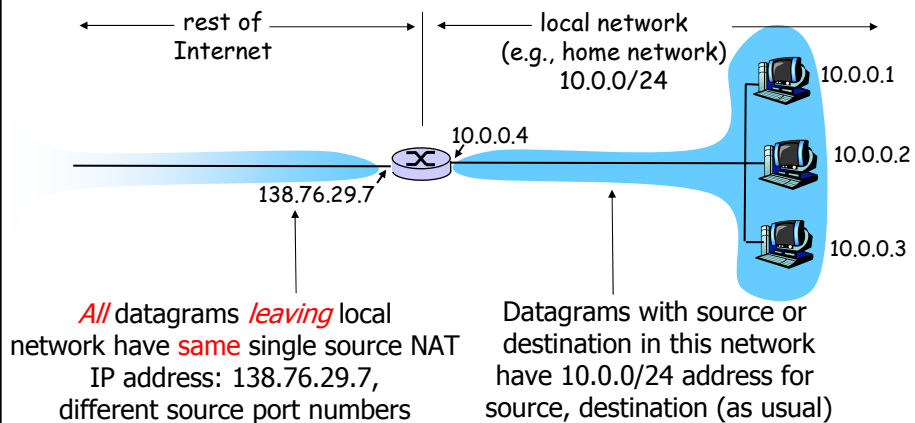
NAT: Network address translation

Motivation: Local network uses just one IP address as far as outside world is concerned:

- Just one IP address for all devices
- Not needed range of addresses from ISP

18

NAT: Network address translation (2.)



19

NAT: Network address translation (3.)

Motivation: Local network uses just one IP address as far as outside world is concerned:

- Range of addresses not needed from ISP: just one IP address for all devices
- Can change addresses of devices in local network without notifying outside world
- Can change ISP without changing addresses of devices in local network
- Devices inside local net not explicitly addressable, visible by outside world (a security plus).

20

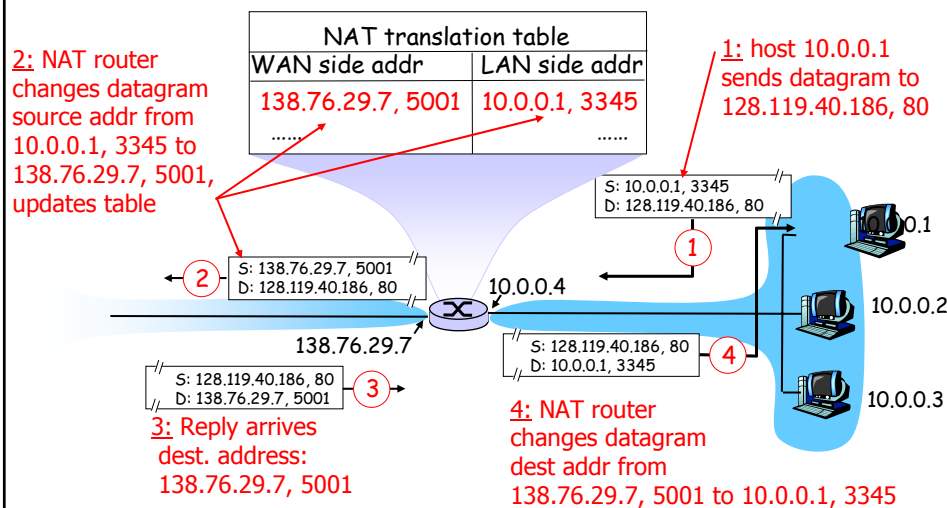
NAT: Network address translation (4.)

Implementation: NAT router must:

- *Outgoing datagrams: Replace* (source IP address, port #) of every outgoing datagram to (NAT IP address, new port #)
 - . . . remote clients/servers will respond using (NAT IP address, new port #) as destination addr.
- *Remember (in NAT translation table)* every (source IP address, port #) to (NAT IP address, new port #) translation pair
- *Incoming datagrams: Replace* (NAT IP address, new port #) in dest fields of every incoming datagram with corresponding (source IP address, port #) stored in NAT table

21

NAT: Network address translation (5.)



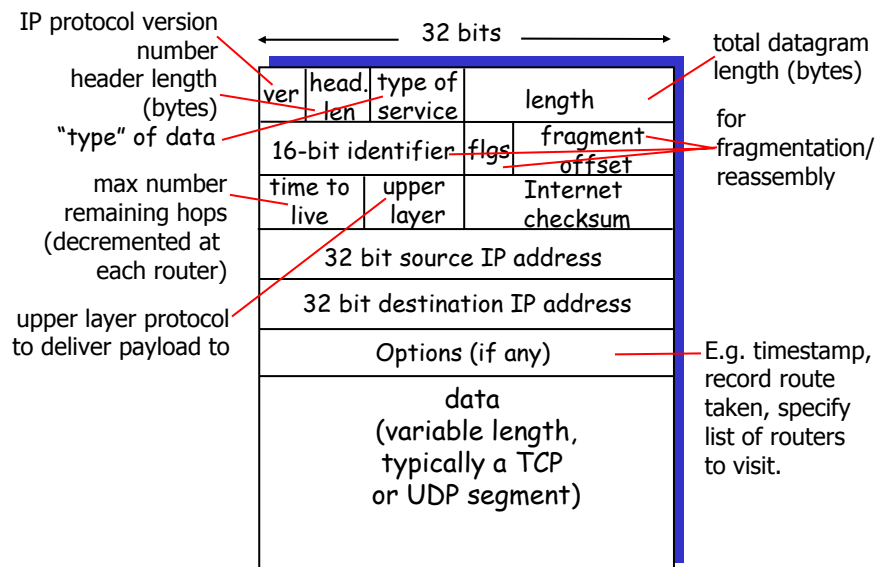
22

NAT: Network address translation (6.)

- 16-bit port-number field:
 - 60,000 simultaneous connections with a single LAN-side address!
- NAT is controversial:
 - Routers should only process up to layer 3
 - Violates end-to-end argument
 - NAT possibility must be taken into account by app designers, e.g., P2P applications
 - Address shortage should instead be solved by IPv6

23

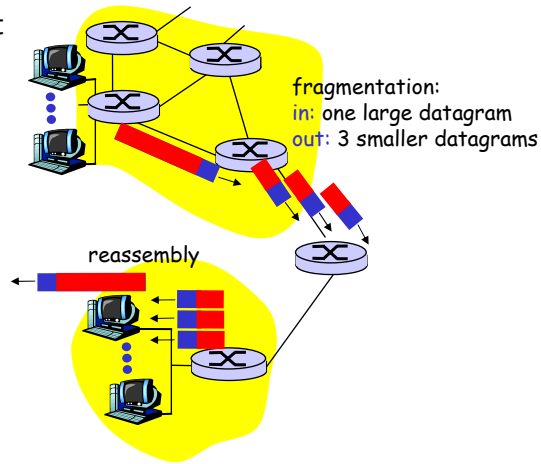
IPv4 Datagram Format



24

IP fragmentation and reassembly

- Network links have MTU (max. transfer size) – largest possible link-level frame.
 - Different link types, different MTUs
- Large IP datagram divided (“fragmented”) within net
 - One datagram becomes several datagrams
 - “Reassembled” only at final destination
 - IP header bits used to identify, order related fragments



25

IP fragmentation and reassembly (2.)

Example

- 4000 byte datagram
- MTU = 1500 bytes

| length | ID | fragflag | offset |
|--------|----|----------|--------|
| =4000 | =x | =0 | =0 |

One large datagram becomes several smaller datagrams

1480 bytes in data field

offset = 1480/8

| length | ID | fragflag | offset |
|--------|----|----------|--------|
| =1500 | =x | =1 | =0 |
| =1500 | =x | =1 | =185 |
| =1040 | =x | =0 | =370 |

26

ICMP: Internet control message protocol

- ❑ Used by hosts, routers, gateways to communication network-level information
 - Error reporting:
 - unreachable host,
 - network, port, protocol
 - Echo request/reply (used by ping)
 - ❑ Network-layer "above" IP:
 - ICMP msgs carried in IP datagrams
 - ❑ **ICMP message:** Type, code plus first 8 bytes of IP datagram causing error
- | Type | Code | description |
|------|------|---|
| 0 | 0 | echo reply (ping) |
| 3 | 0 | dest. network unreachable |
| 3 | 1 | dest host unreachable |
| 3 | 2 | dest protocol unreachable |
| 3 | 3 | dest port unreachable |
| 3 | 6 | dest network unknown |
| 3 | 7 | dest host unknown |
| 4 | 0 | source quench (congestion control – not used) |
| 8 | 0 | echo request (ping) |
| 9 | 0 | route advertisement |
| 10 | 0 | router discovery |
| 11 | 0 | TTL expired |
| 12 | 0 | bad IP header |

27

Address resolution protocol (ARP)

- ❑ Interface between Link layer and Network Layer
- ❑ Allows hosts to query who owns an IP address on the same LAN
- ❑ Owner responds with hardware address
- ❑ Allows changes to link layer to be independent of IP addressing

28

Network layer: Status

- Network layer functions
- IP
- Routing and forwarding
- NAT
- ARP
- IPv6
- Routing