Socket programming

**Goal:** learn how to build client/server application that communicate using sockets

**Socket API**
- Introduced in BSD4.1 UNIX, 1981
- Explicitly created, used, released by apps
- Client/server paradigm
- Two types of transport service via socket API:
  - Unreliable datagram
  - Reliable, byte stream-oriented

**Socket-programming using TCP**

**Socket:** a door between application process and end-end-transport protocol (UDP or TCP)

**TCP service:** reliable transfer of bytes from one process to another
Socket programming with TCP

Client must contact server
- Server process must first be running
- Server must have created socket (door) that welcomes client's contact

Client contacts server by:
- Creating client-local TCP socket
- Specifying IP address, port number of server process

When client creates socket:
- Client TCP establishes connection to server TCP

When contacted by client, server TCP creates new socket for server process to communicate with client
- Allows server to talk with multiple clients

Application viewpoint
TCP provides reliable, in-order transfer of bytes ("pipe") between client and server

Example client-server app:
- Client reads line from standard input (inFromUser stream), sends to server via socket (outToServer stream)
- Server reads line from socket
- Server converts line to uppercase, sends back to client
- Client reads, prints modified line from socket (inFromServer stream)

Input stream: sequence of bytes into process
Output stream: sequence of bytes out of process
Client/server socket interaction: TCP

Server (running on hostid)          Client

create socket,  
port=x, for incoming request:  
welcomeSocket = ServerSocket()

wait for incoming connection request  
connectionSocket = welcomeSocket.accept()

read request from connectionSocket  
write reply to connectionSocket  
close connectionSocket

TCP connection setup

create socket, connect to hostid, port=x  
clientSocket = Socket()

send request using clientSocket  
read reply from connectionSocket  
close clientSocket

Example: Java client (TCP)

```
import java.io.*;
import java.net.*;
class TCPClient {
    public static void main(String argv[]) throws Exception {
        String sentence;
        String modifiedSentence;
        BufferedReader inFromUser =
            new BufferedReader(new InputStreamReader(System.in));
        Socket clientSocket = new Socket("hostname", 6789);
        DataOutputStream outToServer =
            new DataOutputStream(clientSocket.getOutputStream());
```

Create input stream  
Create client socket, connect to server  
Create output stream attached to socket
Example: Java client (TCP), cont.

```java
BufferedReader inFromServer =
  new BufferedReader(new InputStreamReader(clientSocket.getInputStream()));

sentence = inFromUser.readLine();
outToServer.writeBytes(sentence + '\n');
modifiedSentence = inFromServer.readLine();
System.out.println("FROM SERVER: " + modifiedSentence);
clientSocket.close();
```

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Example: Java server (TCP)

```java
import java.io.*;
import java.net.*;

class TCPServer {
  public static void main(String argv[]) throws Exception {
    String clientSentence;
    String capitalizedSentence;
    ServerSocket welcomeSocket = new ServerSocket(6789);

    while(true) {
      Socket connectionSocket = welcomeSocket.accept();
      BufferedReader inFromClient =
        new BufferedReader(new InputStreamReader(connectionSocket.getInputStream()));
```
Example: Java server (TCP), cont

```java
DataOutputStream outToClient = new DataOutputStream(connectionSocket.getOutputStream());
clientSentence = inFromClient.readLine();
capitalizedSentence = clientSentence.toUpperCase() + '\n';
outToClient.writeBytes(capitalizedSentence);
```

End of while loop, loop back and wait for another client connection

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Socket programming with UDP

UDP: no “connection” between client and server

- No handshaking
- Sender explicitly attaches IP address and port of destination
- Server must extract IP address, port of sender from received datagram

UDP: transmitted data may be received out of order, or lost
Client/server socket interaction: UDP

Server (running on hostid)

- Create socket, port=\textit{x}, for incoming request:
  - serverSocket = DatagramSocket()
- Read request from serverSocket
- Write reply to serverSocket specifying client host address, port number

Client

- Create socket, clientSocket = DatagramSocket()
- Create, address (hostid, port=\textit{x}, send datagram request using clientSocket
- Read reply from clientSocket
- Close clientSocket

Example: Java client (UDP)

```java
import java.io.*;
import java.net.*;

class UDPClient {
    public static void main(String args[]) throws Exception {
        BufferedReader inFromUser = new BufferedReader(new InputStreamReader(System.in));
        DatagramSocket clientSocket = new DatagramSocket();
        InetAddress IPAddress = InetAddress.getByName("hostname");
        byte[] sendData = new byte[1024];
        byte[] receiveData = new byte[1024];
        String sentence = inFromUser.readLine();
        sendData = sentence.getBytes();
        String sendAddress = IPAddress.getHostAddress();
        clientSocket.send(new DatagramPacket(sendData, sendData.length, InetAddress.getByName(sendAddress), 1234));
        DatagramPacket reply = new DatagramPacket(receiveData, receiveData.length, clientSocket.getRemoteSocketAddress());
        clientSocket.receive(reply);
        System.out.println("Recieved from server: \" + new String(reply.getData(), 0, reply.getLength()) + "\n");
    }
}
```
Example: Java client (UDP), cont.

Create datagram with data-to-send, length, IP addr, port

DatagramPacket sendPacket = new DatagramPacket(sendData, sendData.length, IPAddress, 9876);
clientSocket.send(sendPacket);

Send datagram to server

DatagramPacket receivePacket = new DatagramPacket(receiveData, receiveData.length);
clientSocket.receive(receivePacket);

Read datagram from server

String modifiedSentence = new String(receivePacket.getData());
System.out.println("FROM SERVER:" + modifiedSentence);
clientSocket.close();

Example: Java server (UDP)

import java.io.*;
import java.net.*;

class UDPServer {
    public static void main(String args[]) throws Exception {
        DatagramSocket serverSocket = new DatagramSocket(9876);
        byte[] receiveData = new byte[1024];
        byte[] sendData = new byte[1024];

        while(true) {
            DatagramPacket receivePacket = new DatagramPacket(receiveData, receiveData.length);
            serverSocket.receive(receivePacket);
        }
    }
}
Example: Java server (UDP), cont

```java
String sentence = new String(receivePacket.getData());
InetAddress IPAddress = receivePacket.getAddress();
int port = receivePacket.getPort();
String capitalizedSentence = sentence.toUpperCase();
sendData = capitalizedSentence.getBytes();
DatagramPacket sendPacket =
    new DatagramPacket(sendData, sendData.length, IPAddress, port);
serverSocket.send(sendPacket);
```

End of while loop, loop back and wait for another client connection

Electronic mail

Three major components:
- User agents
- Mail servers
- Simple mail transfer protocol: smtp

User Agent
- A.k.a. "mail reader"
- Composing, editing, reading mail messages
- E.g., Eudora, pine, elm, Netscape Messenger
- Outgoing, incoming messages stored on server
Electronic mail: mail servers

Mail Servers
- **Mailbox** contains incoming messages (yet to be read) for user
- **Message** queue of outgoing (to be sent) mail messages
- **Smtp protocol** between mail server to send email messages
  - Client: sending mail server
  - “Server”: receiving mail server

Electronic mail: smtp [RFC 821]

- Uses tcp to reliably transfer email msg from client to server, port 25
- Direct transfer: sending server to receiving server
- Three phases of transfer
  - Handshaking (greeting)
  - Transfer
  - Closure
- Command/response interaction
  - **Commands**: ASCII text
  - **Response**: status code and phrase
Sample smtp interaction

S: 220 hamburger.edu
C: HELO crepes.fr
S: 250 Hello crepes.fr, pleased to meet you
C: MAIL FROM: <alice@crepes.fr>
S: 250 alice@crepes.fr... Sender ok
C: RCPT TO: <bob@hamburger.edu>
S: 250 bob@hamburger.edu ... Recipient ok
C: DATA
S: 354 Enter mail, end with "." on a line by itself
C: Do you like ketchup?
C: How about pickles?
C: .
S: 250 Message accepted for delivery
C: QUIT
S: 221 hamburger.edu closing connection

Smtp: final words

Try smtp interaction for yourself:
- telnet servername 25
- See 220 reply from server
- Enter HELO, MAIL FROM, RCPT TO, DATA, QUIT commands

above lets you send email without using email client (reader)

Comparison with http
- HTTP: pull
- Email: push
- Both have ASCII command / response interaction, status codes
- HTTP: multiple objects in file sent in separate connections
- Smtp: multiple message parts sent in one connection
Mail message format

Smtp: protocol for exchanging email msgs
RFC 822: standard for text message format:
- Header lines, e.g.,
  - To:
  - From:
  - Subject: different from smtp commands!
- Body
  - The "message", ASCII characters only
- Line containing only '.'

Message format: multimedia extensions

- MIME: multimedia mail extension, RFC 2045, 2056
- Additional lines in msg header declare MIME content type
MIME types

**Text**
- Example subtypes: plain, html

**Image**
- Example subtypes: jpeg, gif

**Audio**
- Example subtypes: basic (8-bit mu-law encoded), 32kadpcm (32 kbps coding)

**Video**
- Example subtypes: mpeg, quicktime

**Application**
- Other data that must be processed by reader before “viewable”
- Example subtypes: msword, octet-stream

Mail access protocols

- **SMTP**: delivery/storage to receiver’s server
- **Mail access protocol**: retrieval from server
  - **POP**: Post Office Protocol [RFC 1939]
    - Authorization (agent ↔ server) and download
  - **IMAP**: Internet Mail Access Protocol [RFC 1730]
    - More features (more complex)
    - Manipulation of stored msgs on server
POP3 protocol

Authorization phase
- Client commands:
  - **user**: declare username
  - **pass**: password
- Server responses
  - +OK
  - -ERR

Transaction phase, client:
- **list**: list message numbers
- **retr**: retrieve message by number
- **dele**: delete
- **quit**

```
S: +OK POP3 server ready
C: user alice
S: +OK
C: pass hungry
S: +OK user successfully logged on
C: list
S: 1 498
S: 2 912
S: .
C: retr 1
S: <message 1 contents>
S: .
C: dele 1
C: retr 2
S: <message 1 contents>
S: .
C: dele 2
C: quit
S: +OK POP3 server signing off
```

POP3 (more) and IMAP

More about POP3
- Previous example uses “download and delete” mode.
- Bob cannot re-read e-mail if he changes client
- “Download-and-keep”: copies of messages on different clients
- POP3 is stateless across sessions

IMAP
- Keep all messages in one place: the server
- Allows user to organize messages in folders
- IMAP keeps user state across sessions:
  - names of folders and mappings between message IDs and folder name
FTP: the file transfer protocol

- Transfer file to/from remote host
- Client/server model
  - Client: side that initiates transfer (either to/from remote)
  - Server: remote host
- FTP: RFC 959
- FTP server: port 21

FTP: separate control, data connections

- FTP client contacts ftp server at port 21, specifying TCP as transport protocol
- Two parallel TCP connections opened:
  - Control: exchange commands, responses between client, server.
    - "out of band control"
  - Data: file data to/from server
- FTP server maintains "state": current directory, earlier authentication
### FTP commands, responses

**Sample commands:**
- Sent as ASCII text over control channel
- `USER username`
- `PASS password`
- `LIST` return list of file in current directory
- `RETR filename` retrieves (gets) file
- `STOR filename` stores (puts) file onto remote host

**Sample return codes**
- Status code and phrase (as in http)
- 331 Username OK, password required
- 125 data connection already open; transfer starting
- 425 Can’t open data connection
- 452 Error writing file

### Application architectures

- Client-server
- Peer-to-peer (P2P)
- Hybrid of client-server and P2P
Client-server architecture

Server:
- Always-on host
- Permanent IP address
- Server farms for scaling

Clients:
- Communicate with server
- May be intermittently connected
- May have dynamic IP addresses
- Do not communicate directly with each other

Pure P2P architecture

- No always-on server
- Arbitrary end systems directly communicate
- Peers are intermittently connected and change IP addresses
- Example: Gnutella

Highly scalable but difficult to manage
Generations of P2P

- 1st Generation: Centralized file list
  - Napster
  - He who controls central file is responsible legally
- 2nd Generation: Decentralized file lists
  - Gnutella, FastTrack
  - Improvements – optimizations of decentralized search
- 3rd Generation: No file lists
  - Freenet, WASTE, Entropy, MUTE
  - Anonymity built in

The Good, Bad, and Ugly of P2P

- The Good
  - Security based on social contract
  - Free exchange of ideas
  - Everyone’s computer can contribute to the greater good
- The Bad
  - Avoids most security: Can be used for piracy
  - “The Tragedy of the Commons” – not supported by all IP
Hybrid of client-server and P2P

Skype
- Internet telephony app
- Finding address of remote party: centralized server(s)
- Client-client connection is direct (not through server)

Instant messaging
- Chatting between two users is P2P
- Presence detection/location centralized:
  - User registers its IP address with central server when it comes online
  - User contacts central server to find IP addresses of buddies

P2P file sharing

Example
- Alice runs P2P client application on her notebook computer
- Intermittently connects to Internet; gets new IP address for each connection
- Asks for “Hey Jude”
- Application displays other peers that have copy of Hey Jude.
- Alice chooses one of the peers, Bob.
- File is copied from Bob’s PC to Alice’s notebook: HTTP
- While Alice downloads, other users uploading from Alice.
- Alice’s peer is both a Web client and a transient Web server.

All peers are servers = highly scalable!
**P2P: centralized directory**

Original “Napster” design

1) When peer connects, it informs central server:
   - IP address
   - Content

2) Alice queries for “Hey Jude”

3) Alice requests file from Bob

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**P2P: problems with centralized directory**

- Single point of failure
- Performance bottleneck
- Copyright infringement

File transfer is decentralized, but locating content is highly centralized
Query flooding: Gnutella

- Fully distributed
  - No central server
- Public domain protocol
- Many Gnutella clients implementing protocol

Overlay network: graph

- Edge between peer X and Y if there's a TCP connection
- All active peers and edges are part of the overlay net
- Edge is not a physical link
- Given peer will typically be connected with < 10 overlay neighbors

Gnutella: protocol

- Query message sent over existing TCP connections
- Peers forward Query message
- QueryHit sent over reverse path

File transfer: HTTP

Scalability: limited scope flooding
Gnutella: peer joining

1. Joining peer X must find some other peer in Gnutella network: use list of candidate peers
2. X sequentially attempts to make TCP with peers on list until connection setup with Y
3. X sends Ping message to Y; Y forwards Ping message.
4. All peers receiving Ping message respond with Pong message
5. X receives many Pong messages. It can then setup additional TCP connections

Exploiting heterogeneity: KaZaA

- Each peer is either a group leader or assigned to a group leader.
  - TCP connection between peer and its group leader.
  - TCP connections between some pairs of group leaders.
- Group leader tracks the content in all its children.
KaZaA: querying

- Each file has a hash and a descriptor
- Client sends keyword query to its group leader
- Group leader responds with matches:
  - For each match: metadata, hash, IP address
- If group leader forwards query to other group leaders, they respond with matches
- Client then selects files for downloading
  - HTTP requests using hash as identifier sent to peers holding desired file

KaZaA tricks

- Limitations on simultaneous uploads
- Request queuing
- Incentive priorities
- Parallel downloading

For more info:
- J. Liang, R. Kumar, K. Ross, “Understanding KaZaA,”
  (available via cis.poly.edu/~ross/p2pTheory)
BitTorrent

- Scalable way of downloading bulk data
  - Breaks files into multiple data blocks
  - After downloading a data block, clients help upload it to other clients
- Many legitimate uses
  - Software distribution
    - Unix distributions, e.g., Linux
    - World of Warcraft
  - Media distribution
    - bittorrent.com (commercial service)

BitTorrent terms

- Client
  - Leecher (Sauger): initial state
  - Seeder: once it has the full dataset (file)
- Tracker
  - Coordinator
- Torrent file
  - Meta information
- Torrent
  - The overlay
**BitTorrent – joining a torrent**

1. Obtain the metadata file
2. Contact the tracker
3. Obtain a peer list (contains seeds & leechers)
4. Contact peers from that list for data

Peers divided into:
- **seeders**: have the entire file
- **leechers**: still downloading

**BitTorrent – exchanging data**

- Verify pieces using hashes
- Download sub-pieces in parallel
- Advertise received pieces to the entire peer set (i.e., the connected peers of a single peer)
  - HAVE Message
- Look for the rarest pieces
BitTorrent – philosophy

- Philosophy
  - Based on Tit-for-tat
  - Incentive - Uploading while downloading

- Choking algorithm
  - Controls whom you are sending data
  - Choke – disable sending of data to peer
  - Unchoke – enable sending of data to peer
  - Preferred peers – those that are sending you data
  - Optimistic unchoke – to enable download from peers
  - Snubbing – State resulting in bad performance

BitTorrent – unchoking

- Periodically calculate data-receiving rates
- Upload to (unchoke) the fastest downloaders

- Optimistic unchoking
  - Periodically select a peer at random and upload to it
  - Continuously look for the fastest partners
Overall Architecture

Web page with link to .torrent

Web Server

Tracker

Peer [Leech]
Downloader “US”

Peer [Seed]

Peer [Leech]

Web page with link to .torrent

Get-announce

Peer [Leech]

Peer [Seed]
Overall Architecture

Web page with link to .torrent

Response-peer list

Peer [Leech] Downloader “US”

Tracker

Peers

B

A

C

Web Server

Shake-hand

Peer [Leech]

Peer [Seed]
Overall Architecture

Web page with link to .torrent

A

Peer [Leech]
Downloader "US"

B

Peer [Leech]

C

Peer [Seed]

Web Server
Tracker

pieces

Overall Architecture

Web page with link to .torrent

A

Peer [Leech]
Downloader "US"

B

Peer [Leech]

C

Peer [Seed]

Web Server
Tracker

pieces

pieces

pieces

pieces
Chat systems

- Popular
- Highly interactive services
- Allow to exchange text messages interactively with other people around the world

Types
- IRC (Internet relay chat)
- Web-Chat
  - HTML based
  - Applet based
- Instant messagengers (ICQ, AIM, MIM, ...)
- Others
**IRC**

- Wideley used – relatively old
  - 2003: 5 biggest networks > 250,000 users
- Channels
- User = unique nickname
- Commands: PRIVMSG, JOIN, ISON, NICK
  - User input: /MSG nick2 Hi
  - IRC Format: :nick1 PRIVMSG nick2 :Hi!
  - Output: *nick1* Hi!

**Properties of IRC**

- IRC operators administer IRC network
- Port 6667
- Well-defined protocol (RFC 1459)
- Client/server to server network
  - Allows direct computer-to-computer connections (DCC)
- Usage
  - User connects to public IRC server
  - Joins channel
  - Chats with others
  - Shares files via DCC
Sample IRC network

IRC: Dangers

- Means of spreading worms
- Susceptible to flooding
- Can be embedded in trojans and act as hostile server unnoticed
Malicious IRC usages

- Exchange exploit information
  - Passwords, etc. ...
- Common protocol for communication between attackers
- Use IRC as control channel
  - Use private IRC networks
- Distributed Denial of Service attacks
  - Clone/Flood/War bots simulate multiple users in a channel
  - Bots spread and infect lots of computers that log into the channel
  - Attacker sends a command through IRC causing all bots to simultaneously flood packets to a target, e.g., Yahoo, eTrade, Amazon.com, ...)

Web-Chat

- Widely used – newer
- Simple user interface
- Lots of systems using different protocols
- HTML based
  - Interface: browser
  - Protocol: HTTP
- Applet based
  - Interface: applet window
  - Protocol: customer or IRC
Application-layer: Summary

Most importantly: learned about protocols

- Typical request/reply message exchange:
  - Client requests info or service
  - Server responds with data, status code

- Message formats:
  - Headers: fields giving info about data
  - Data: info being communicated

- Control vs. data msgs
  - In-band, out-of-band
- Centralized vs. decentralized
- Stateless vs. stateful
- Reliable vs. unreliable msg transfer
- “Complexity at network edge”