Application Layer

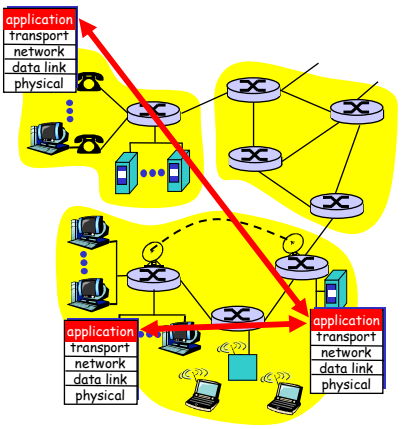
Goals:

- Conceptual aspects of network application protocols
  - Client server paradigm
  - Service models
- Review protocols by examining popular application-level protocols
  - HTTP
  - DNS

Applications and application-layer protocols

Application: communicating, distributed processes
- Running in network hosts in "user space"
- Exchange messages to implement app
- E.g., email, file transfer, the Web

Application-layer protocols
- One "piece" of an app
- Define messages exchanged by apps and actions taken
- User services provided by lower layer protocols



Client-server paradigm

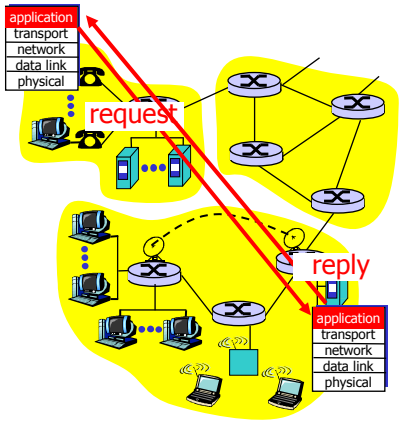
Typical network app has two pieces: *client* and *server*

Client:

- Initiates contact with server ("speaks first")
- Typically requests service from server,
- E.g., request WWW page, send email

Server:

- Provides requested service to client
- E.g., sends requested WWW page, receives/stores received email



3

Services provided by Internet transport protocols

TCP service:

- *Connection-oriented:* setup required between client, server
- *Reliable transport* between sending and receiving process
- *Flow control:* sender won't overwhelm receiver
- *Congestion control:* throttle sender when network overloaded
- *Does not providing:* timing, minimum bandwidth guarantees

UDP service:

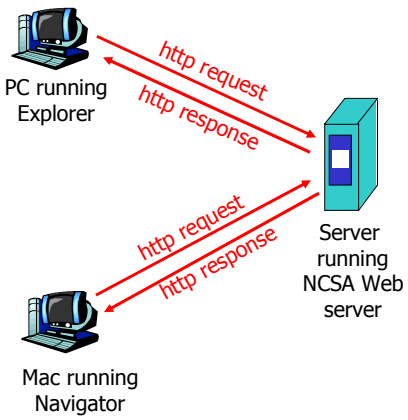
- Unreliable data transfer between sending and receiving process
- Does not provide: connection setup, reliability, flow control, congestion control, timing, or bandwidth guarantee

4

WWW: The HTTP protocol

HTTP: hypertext transfer protocol

- WWW's application layer protocol
- Client/server model
  - *Client:* browser that requests, receives, "displays" WWW objects
  - *Server:* WWW server sends objects in response to requests



5

HTTP - timeline

- Mar 1990 CERN labs document proposing Web
- Jan 1992 HTTP/0.9 specification
- Dec 1992 Proposal to add MIME to HTTP
- Feb 1993 UDI (Universal Document Identifier) Network
- Mar 1993 HTTP/1.0 first draft
- Jun 1993 HTML (1.0 Specification)
- Oct 1993 URL specification
- Nov 1993 HTTP/1.0 second draft
- Mar 1994 URI in WWW
- May 1996 HTTP/1.0 Informational, RFC 1945
- Jan 1997 HTTP/1.1 Proposed Standard, RFC 2068
- Jun 1999 HTTP/1.1 Draft Standard, RFC 2616
- 2001 HTTP/1.1 Formal Standard
- …

6

The HTTP protocol: More

HTTP: TCP transport service

- Client initiates TCP connection (creates socket) to server, port 80
- Server accepts TCP connection from client
- http messages (application-layer protocol messages) exchanged between browser (http client) and WWW server (http server)
- TCP connection closed

HTTP is "stateless"

- Server maintains no information about past client requests

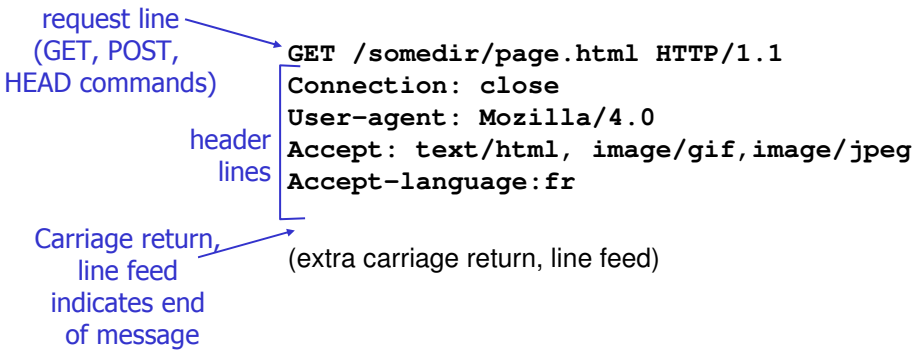
aside

Protocols that maintain "state" are complex!

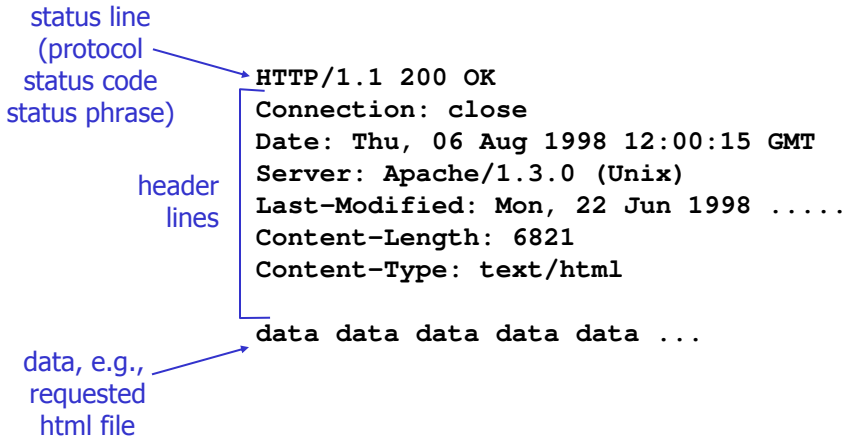
- Past history (state) must be maintained
- If server/client crashes, their views of "state" may be inconsistent, must be reconciled

HTTP message format: Request

- Two types of http messages: *Request, response*
- http request message:
  - ASCII (human-readable format)



HTTP message format: Reply



HTTP reply status codes

In first line in server → client response message.

A few sample codes:

200 OK

- Request succeeded, requested object later in this message

301 Moved Permanently

- Requested object moved, new location specified later in this message (Location:)

400 Bad Request

- Request message not understood by server

404 Not Found

- Requested document not found on this server

505 HTTP Version Not Supported

HTTP request methods

- Methods
  - GET
  - HEAD
  - POST
  - PUT
  - Delete

11

The HTTP protocol: Even more

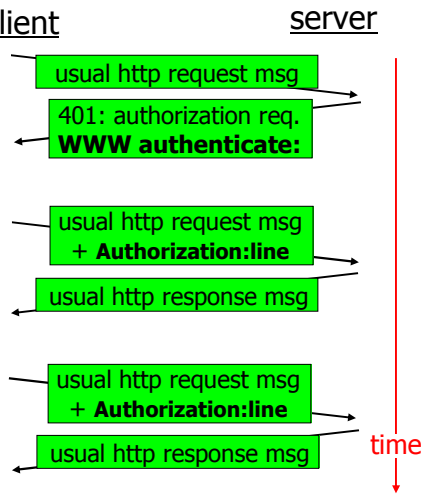
- Non-persistent connection:
  One object in each TCP connection
  - Some browsers create multiple TCP connections *simultaneously* – one per object
- Persistent connection:
  Multiple objects transferred within one TCP connection
- Pipelined persistent connections:
  Multiple requests issued without waiting for response

12

User-server interaction: Authentication

Authentication goal: Control access to server documents

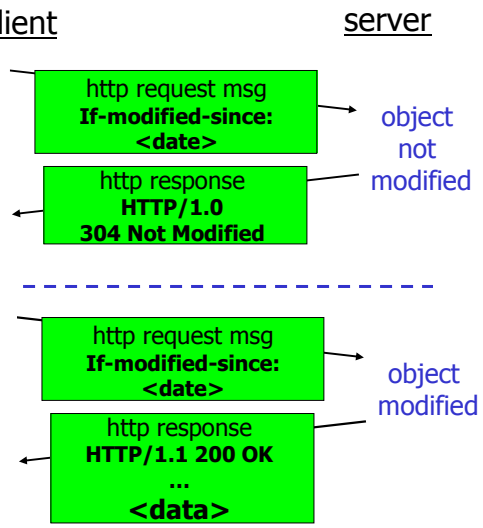
- Stateless: Client must present authorization in each request
- Authorization: Typically name, password
  - `Authorization:` header line in request
  - If no authorization, server refuses access, sends `WWW authenticate:` header line in response



13

User-server interaction: Conditional GET

- Goal: Don't send object if client has up-to-date stored (cached) version
- Client: Specify date of cached copy in http request
  `If-modified-since: <date>`
- Server: Response contains no object if cached copy up-to-date:
  `HTTP/1.0 304 Not Modified`



14

User-server state: Cookies

Most Web sites use cookies

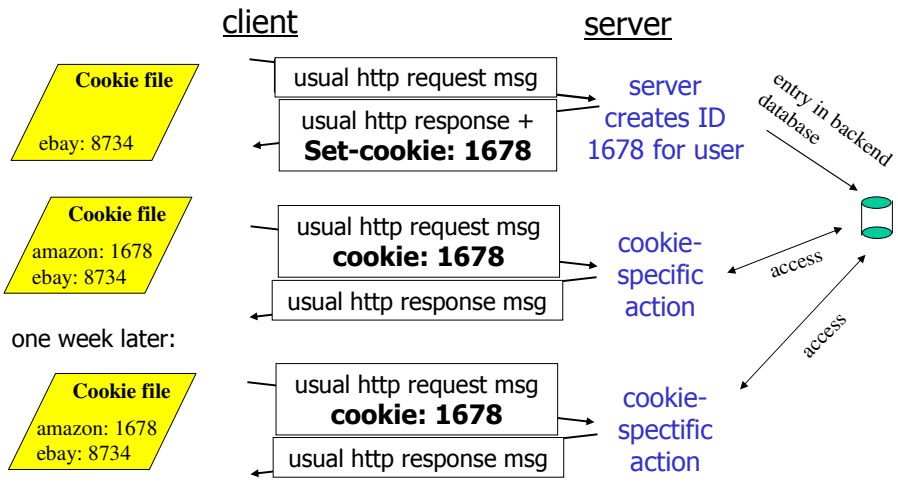
Four components:

1) Cookie header line of HTTP *response* message
2) Cookie header line in HTTP *request* message
3) Cookie file kept on user's host, managed by user's browser
4) Back-end database at Web site

Example:

- Susan access Internet always from same PC
- She visits a specific e-commerce site for first time
- When initial HTTP requests arrives at site, site creates a unique ID and creates an entry in backend database for ID

Cookies: Keeping "state"



Cookies: Keeping "state" (2)

What cookies can bring:

- Authorization
- Shopping carts
- Recommendations
- User session state (Web e-mail)

aside

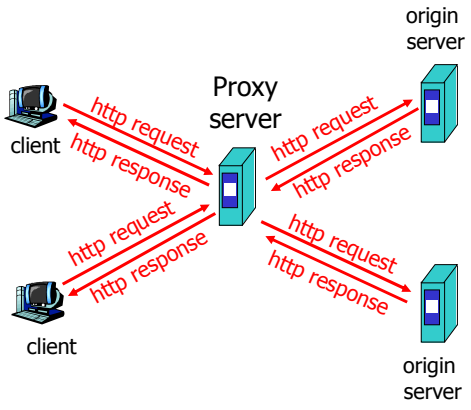
Cookies and privacy:

- Cookies permit sites to learn a lot about you
- You may supply name and e-mail to sites
- Search engines use redirection & cookies to learn yet more
- Advertising companies obtain info across sites

Web caches (proxy server)

Goal: satisfy client request without involving origin server

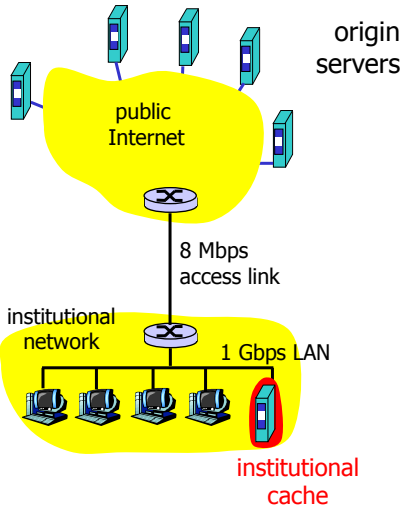
- User sets browser: WWW accesses via web cache
- Client sends all http requests to web cache
  - If object at web cache, web cache immediately returns object in http response
  - Else requests object from origin server, then returns http response to client



Why WWW caching?

Assume: Cache is "close" to client (e.g., in same network)

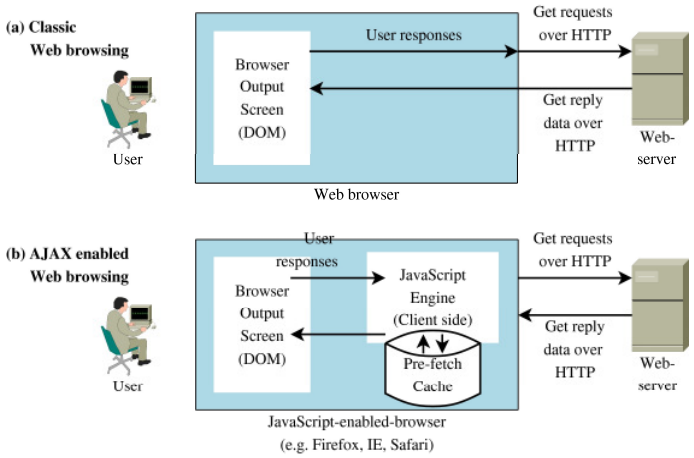
- Smaller response time: cache "closer" to client
- Decrease traffic to distant servers
  - Link out of institutional/local ISP network often bottleneck



19

Web 2.0: E.g., AJAX enabled apps

- E.g.: Google Maps: a canonical AJAX application



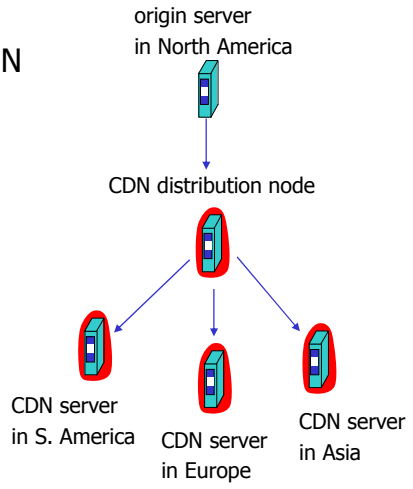
20

Content distribution networks (CDNs)

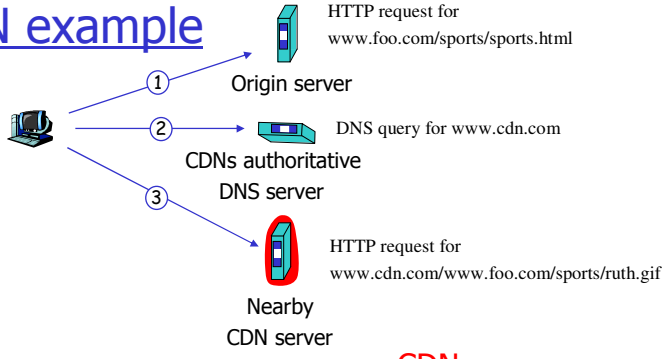
Content providers are the CDN customers.

Content replication

- CDN company installs hundreds of CDN servers throughout Internet
  - In lower-tier ISPs, close to users
- CDN replicates its customers' content in CDN servers. When provider updates content, CDN updates servers



CDN example



Origin server

- www.foo.com
- Distributes HTML
- Replaces:
  http://www.foo.com/sports.ruth.gif
  with
  http://www.cdn.com/www.foo.com/sports/ruth.gif

CDN company

- cdn.com
- Distributes gif files
- Uses its authoritative DNS server to route redirect requests

More about CDNs

Routing requests

- CDN creates a "map", indicating distances from leaf ISPs and CDN nodes
- When query arrives at authoritative DNS server
  - Server determines ISP from which query originates
  - Uses "map" to determine best CDN server

Not just Web pages

- Streaming stored audio/video
- Streaming real-time audio/video
  - CDN nodes create application-layer overlay network

23

DNS: Domain Name System

People: many identifiers:

- SSN, name, Passport #

Internet hosts, routers:

- IP address (32 bit) – used for addressing datagrams
- "name", e.g., gaia.cs.umass.edu – used by humans

Q: Map between IP addresses and name?

- Secure Domain Name System (DNS)
  Dynamic Update: RFC 3007

24

DNS: Domain Name System

Domain Name System:

- *Distributed database:* Implemented in hierarchy of many *name servers*
- *Application-layer protocol:* Host, routers, name servers communicate to *resolve* names (address/name translation)
  - Core Internet function implemented as application-layer protocol
  - Complexity at network's "edge"

25

DNS name servers

Why not centralize DNS?

- Single point of failure
- Traffic volume
- Distant centralized database
- Maintenance

Does not *scale!*

26

DNS name servers (2)

No server has all name-to-IP address mappings

Local name servers:

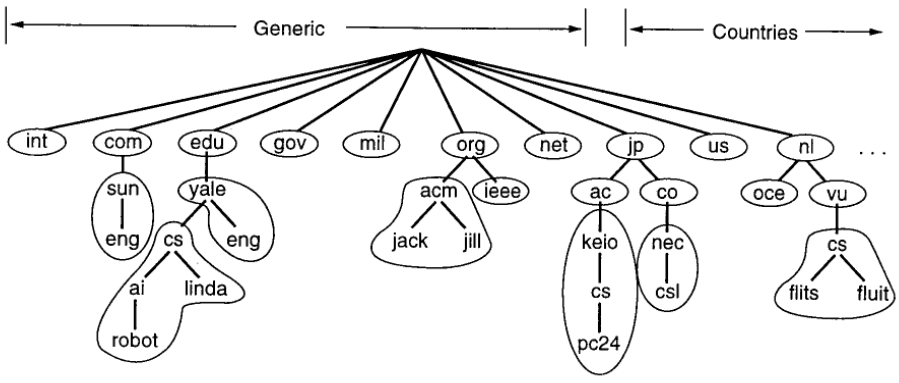
- Each ISP, company has *local (default) name server*
- Host DNS query first goes to local name server

Authoritative name server:

- For a host: stores that host's IP address, name
- Can perform name/address translation for that host's name

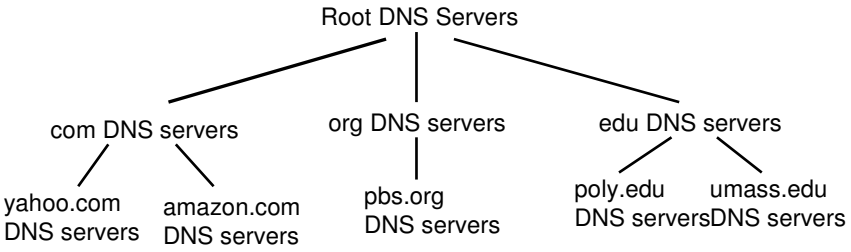
27

DNS: Hierarchical naming tree



28

Distributed, hierarchical database



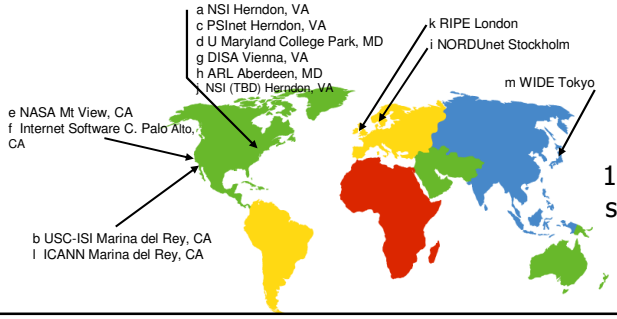
Client wants IP for www.amazon.com; 1st approx:

- Client queries a root server to find com DNS server
- Client queries com DNS server to get amazon.com DNS server
- Client queries amazon.com DNS server to get IP address for www.amazon.com

29

DNS: Root name servers

- Contacted by local name server that can not resolve name
- Root name server:
  - Contacts authoritative name server if name mapping not known
  - Gets mapping
  - Returns mapping to local name server
  - Some use anycast



13 root name servers worldwide

30

TLD and authoritative servers

- Top-level domain (TLD) servers: Responsible for com, org, net, edu, etc, and all top-level country domains uk, fr, ca, jp.
  - Network solutions maintains servers for com TLD
  - Educause for edu TLD
- Authoritative DNS servers: Organization's DNS servers, providing authoritative hostname to IP mappings for organization's servers (e.g., Web and mail).
  - Can be maintained by organization or service provider

31

Local name server

- Does not strictly belong to hierarchy
- Each ISP (residential ISP, company, university) has one.
  - Also called "default name server"
- When a host makes a DNS query, query is sent to its local DNS server
  - Acts as a proxy, forwards query into hierarchy.

32

DNS records

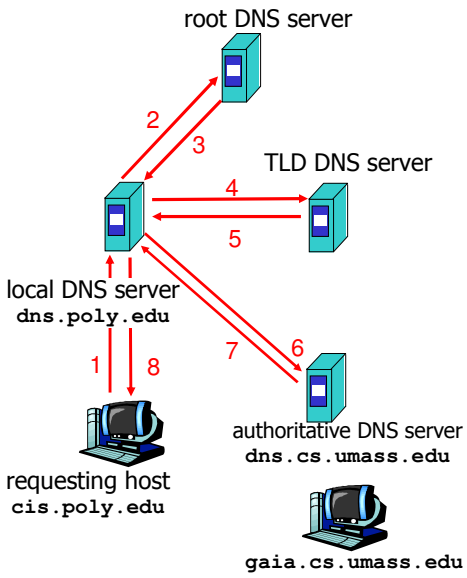
DNS: Distributed db storing resource records (RR)

RR format: `(name, value, type, ttl)`

- Type=A
  - **`name`** is hostname
  - **`value`** is IP address
- Type=NS
  - **`name`** is domain (e.g., foo.com)
  - **`value`** is IP address of authoritative name server for this domain
- Type=CNAME
  - for alias
- Type=MX
  - for mail

Example

- Host at cis.poly.edu wants IP address for gaia.cs.umass.edu



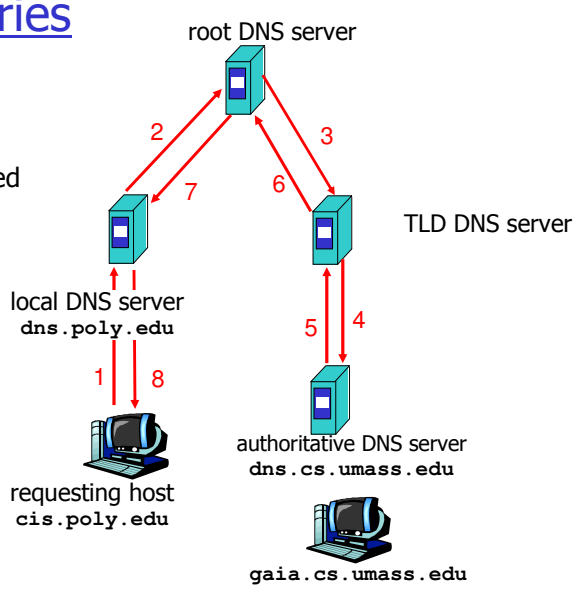
Recursive queries

Recursive query:

- Puts burden of name resolution on contacted name server
- Heavy load?

Iterated query:

- Contacted server replies with name of server to contact
- "I don't know this name, but ask this server"



35

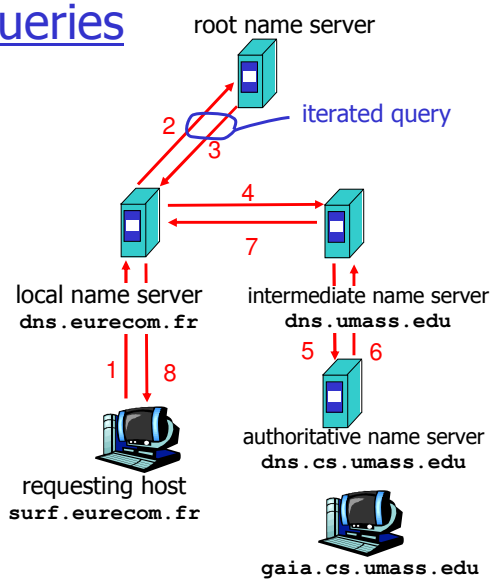
DNS: Iterative queries

Recursive query:

- Puts burden of name resolution on contacted name server
- Heavy load?

Iterated query:

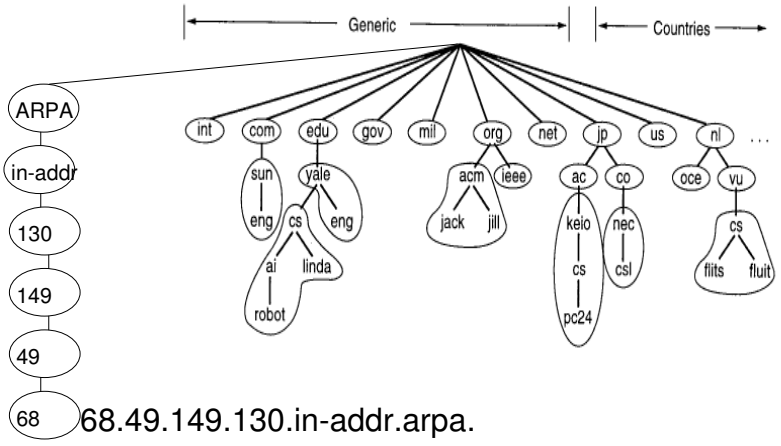
- Contacted server replies with name of server to contact
- "I don't know this name, but ask this server"



36

Mapping IP address to names

- Special domain: ARPA



68.49.149.130.in-addr.arpa.

37