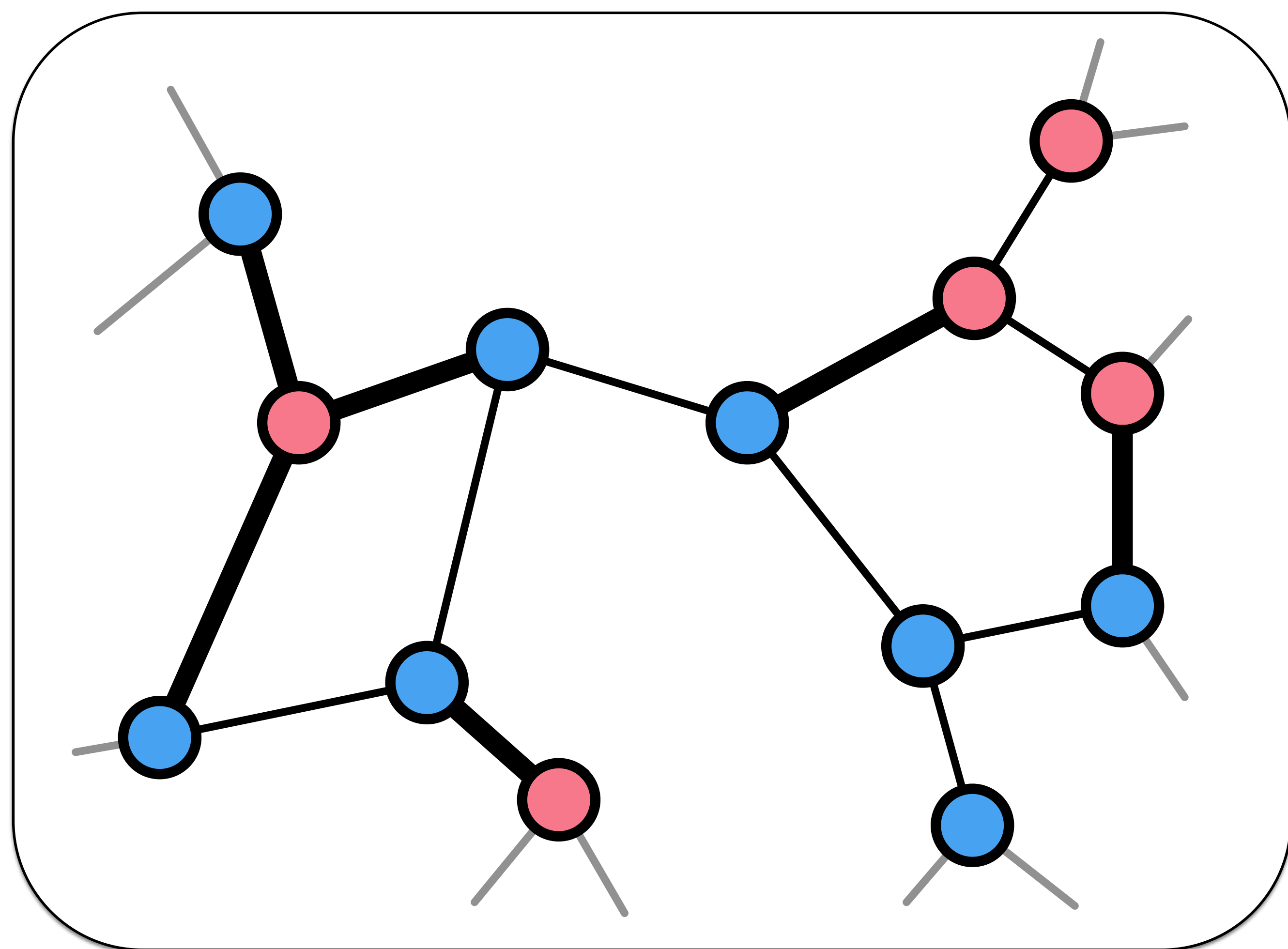


Large Cuts with Local Algorithms

Juho Hirvonen
Joel Rybicki
Jukka Suomela
Aalto University

Online manuscript
arXiv:1402.2542

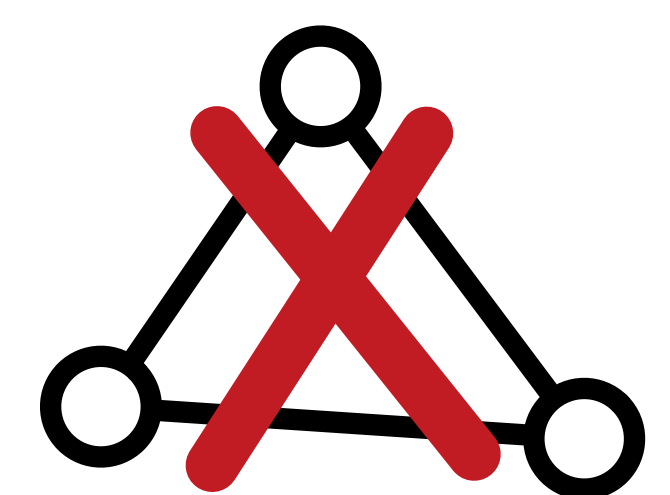
Stefan Schmid
TU Berlin & T-Labs



$$c: V \rightarrow \{\text{blue}, \text{red}\}$$

Input graph

- ◆ d -regular
- ◆ triangle-free

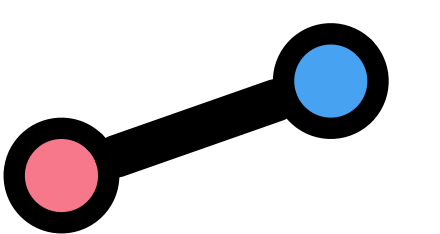


Model of computing

- ◆ n nodes = processors
- ◆ m edges = communication links

Finding large cuts

- ◆ maximise number of *cut edges*



Main result

- ◆ *one-round* randomized algorithm
- ◆ expected number of cut edges is

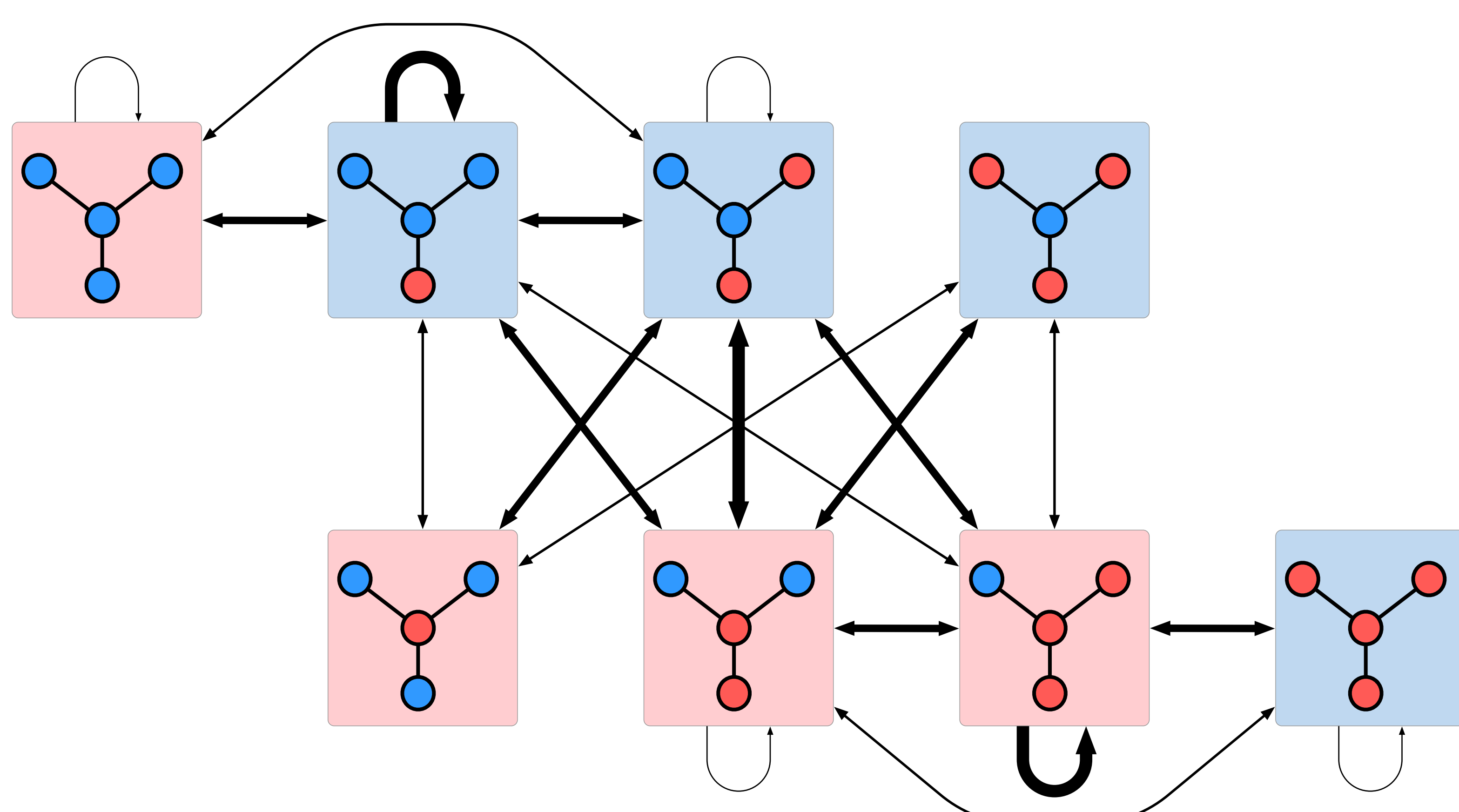
$$\left(\frac{1}{2} + \frac{0.23125}{\sqrt{d}} \right) m$$

Lower bound

- ◆ no sub-logarithmic time algorithm can produce a cut larger than

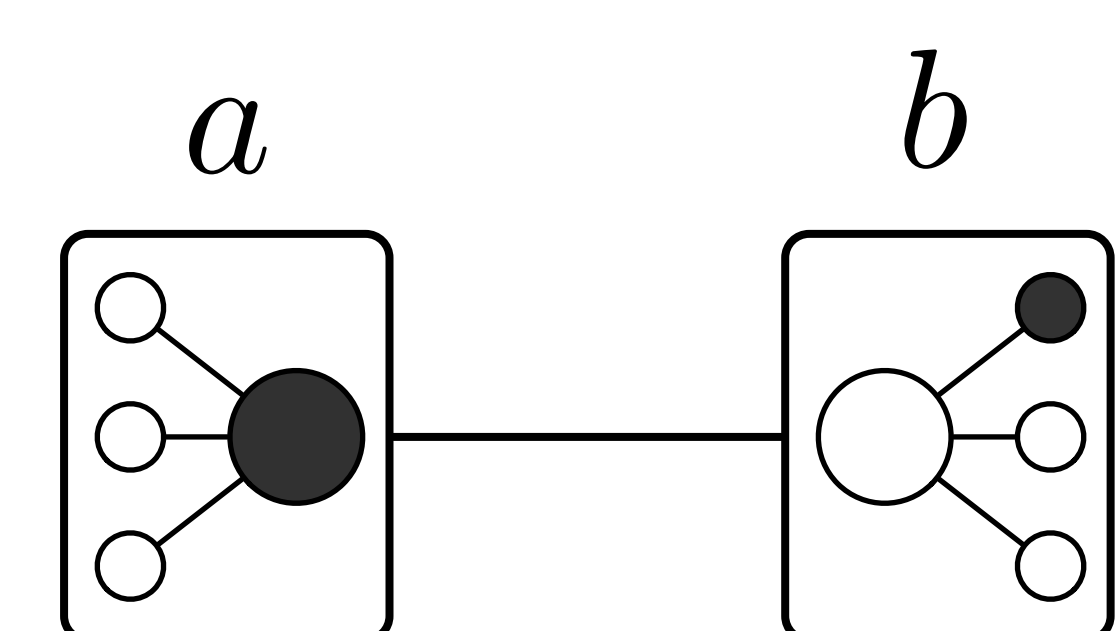
$$\left(\frac{1}{2} + \frac{\sqrt{d-1}}{d} \right) m$$

Computational Algorithm Design



Neighbourhood graphs

heavy cut = good local algorithm



$$(x_a \vee x_b) : w(a, b)$$

$$(\bar{x}_a \vee \bar{x}_b) : w(a, b)$$

Weighted MAX-SAT

find *optimal* algorithms for any fixed d