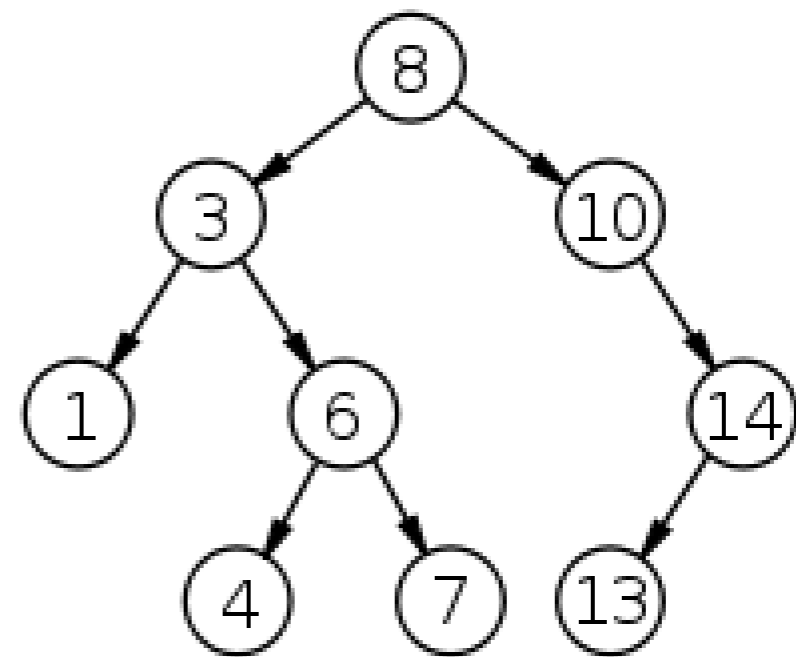


SplayNets

Stefan Schmid, Chen Avin, Christian Scheideler,
Bernhard Häupler, Zvi Lotker

Splay Trees

- Self-adjusting **Binary Search Tree** [ST85]
- A sequence of lookup requests
- “**splay operation**” -
move requested node to the root
- Frequent nodes get close to the root
- Competitive against optimal static tree

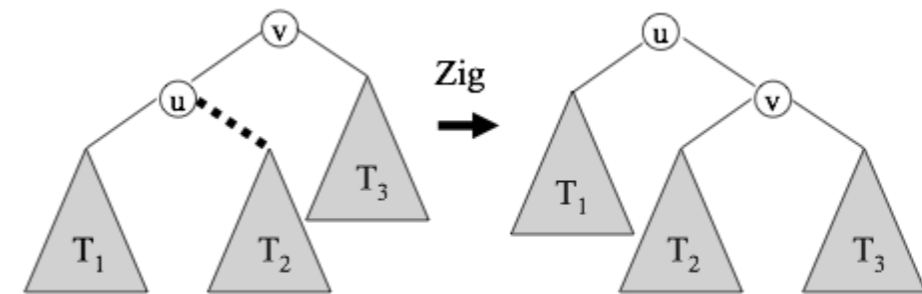


Splay Trees

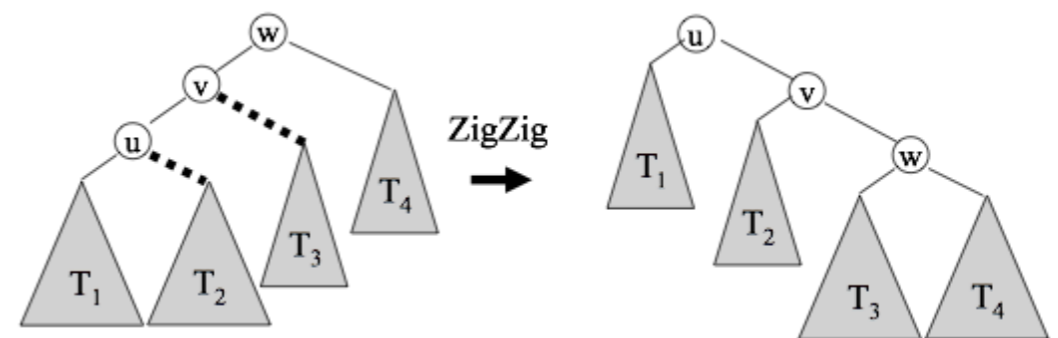
- Self-adjusting **Binary Search Tree** [ST85]

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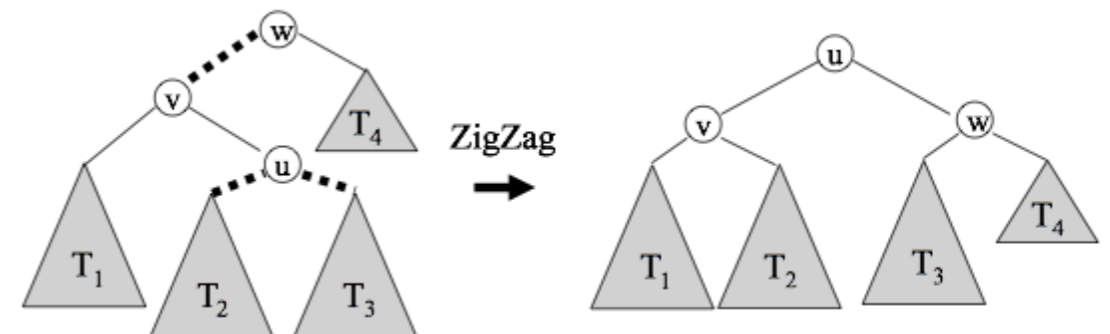
- “**splay operation**” -
move requested node
to the root



- Frequent nodes get
close to the root



- Competitive against
optimal static tree



Our Model

- Instead of a set of requests starting at the root, we have a sequence of **request pairs**

$$\sigma = (\sigma_0, \sigma_1 \dots \sigma_{m-1}) \quad \sigma_t = (u, v) \in V \times V$$

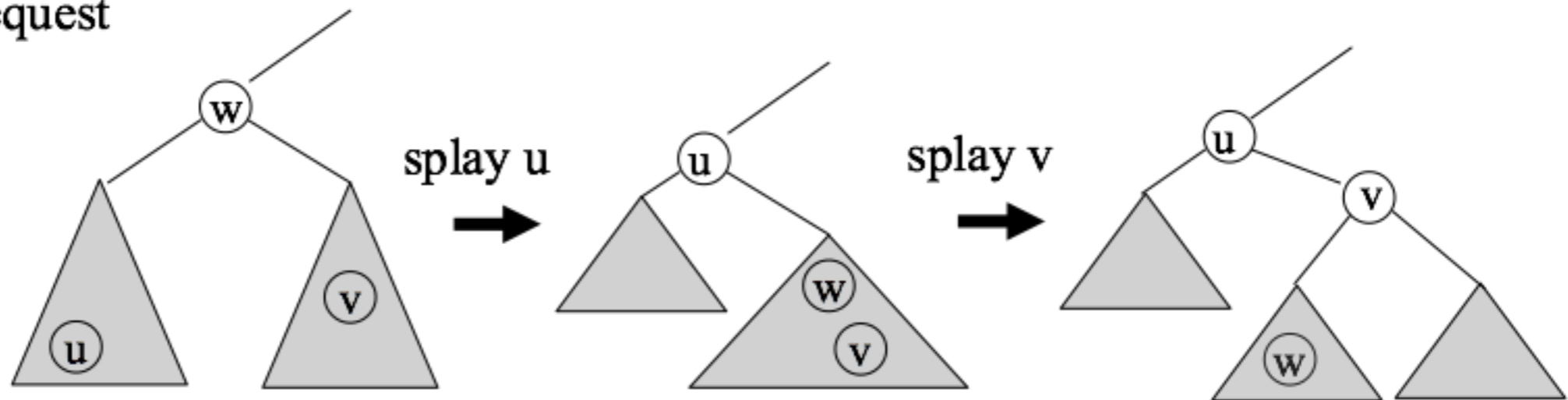
- **What is an optimal static search tree?**
- Can be computed in poly time using dynamic programming.
- But that doesn't provide much insight into problem.

Idea: Double Splay

Algorithm 2 Double Splay Algorithm DS

- 1: (* upon request (u, v) in T *)
 - 2: $w := \alpha_T(u, v)$ \longleftarrow Least common ancestor
 - 3: $T' := \mathbf{splay}$ u to root of $T(w)$
 - 4: \mathbf{splay} v to root of $T'(u)$
-

upon request
 (u, v) :



Analysis: extend Entropy

- **Entropy:** $X \sim p(x), \{x_1, \dots, x_n\}$

$$H(X) = \sum_{i=1}^n p(x_i) \log_2 \frac{1}{p(x_i)} \leq \log n$$

- **Empirical Entropy:** $\hat{X}(\sigma) = \{f(x_1), \dots, f(x_n)\}$

$$H(\hat{X}), H(\hat{Y}), H(\hat{X}, \hat{Y}), H(\hat{X}|\hat{Y})$$

Our Results

Theorem: Let σ be an arbitrary sequence of communication requests, then for any initial BST T_0 ,

$$\text{Cost}(\text{DS}, T_0, \sigma) = O(H(\hat{X}) + H(\hat{Y}))$$

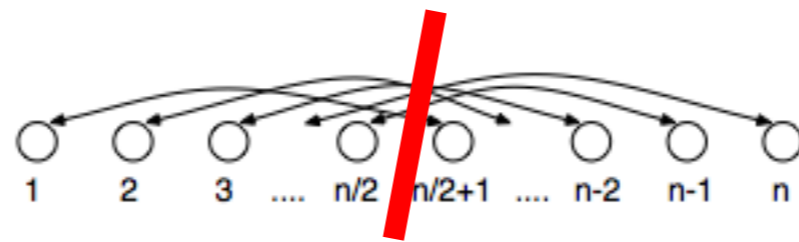
where $H(\hat{X})$ and $H(\hat{Y})$ are the empirical entropies of the sources and the destinations in σ , respectively. Moreover, for any optimal BST network T ,

$$\text{Cost}(T, \sigma) = \Omega(H(\hat{Y} | \hat{X}) + H(\hat{X} | \hat{Y})).$$

Further Results

- A cut based lower bound: there is a cut C

$$\text{Cost}(T, \sigma) = \Omega(H(C(\sigma)))$$



- An expansion lower bound
- Optimality in special cases:
 - Requests form product distribution, rooted binary tree, laminated sets, ...

Thank you!