

Towards Jamming-Resistant and Competitive Medium Access in the SINR Model

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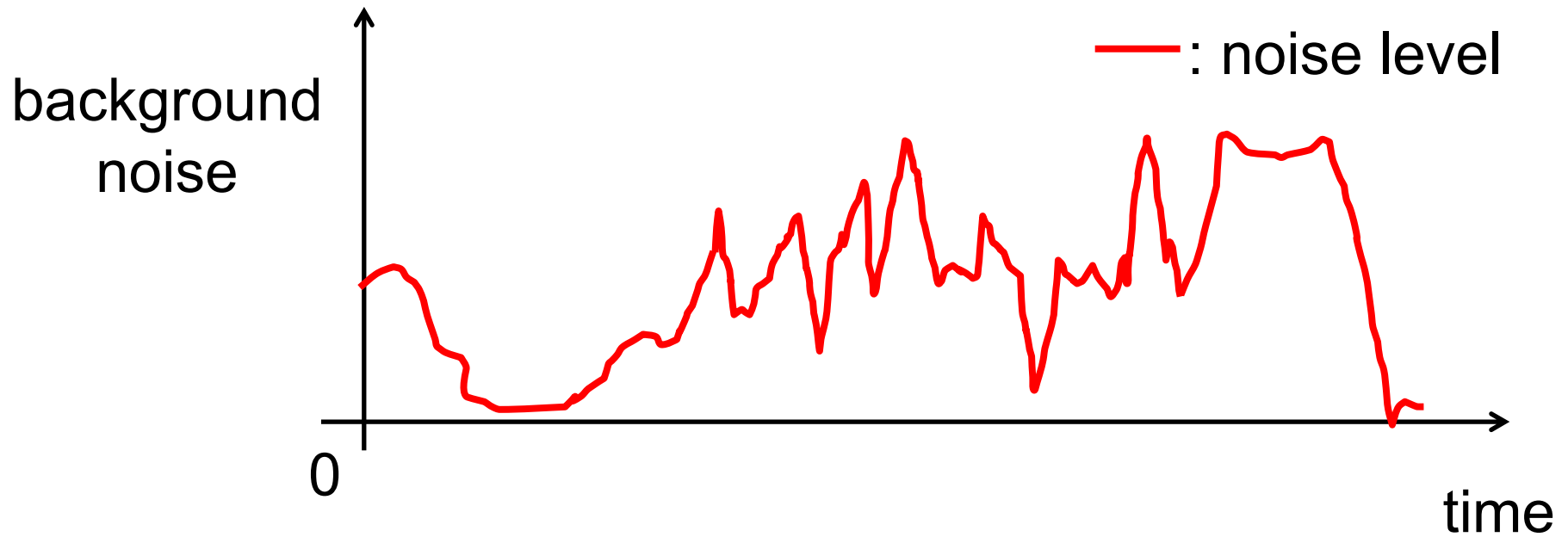
Motivation

Channel availability hard to model:

- Background noise
- Temporary obstacles
- Mobility
- Co-existing networks
- Jammer

Motivation

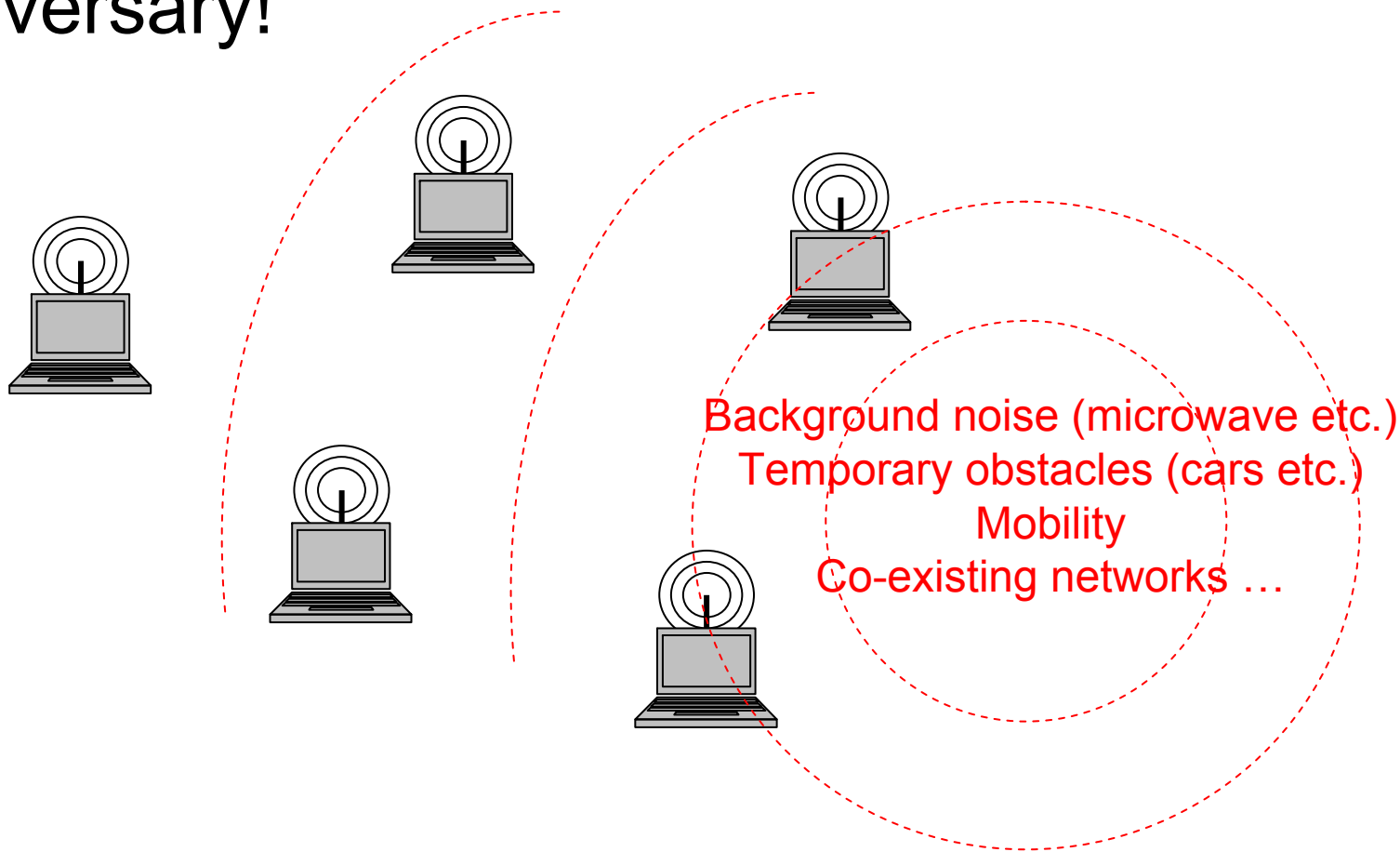
Real world:



How to model this???

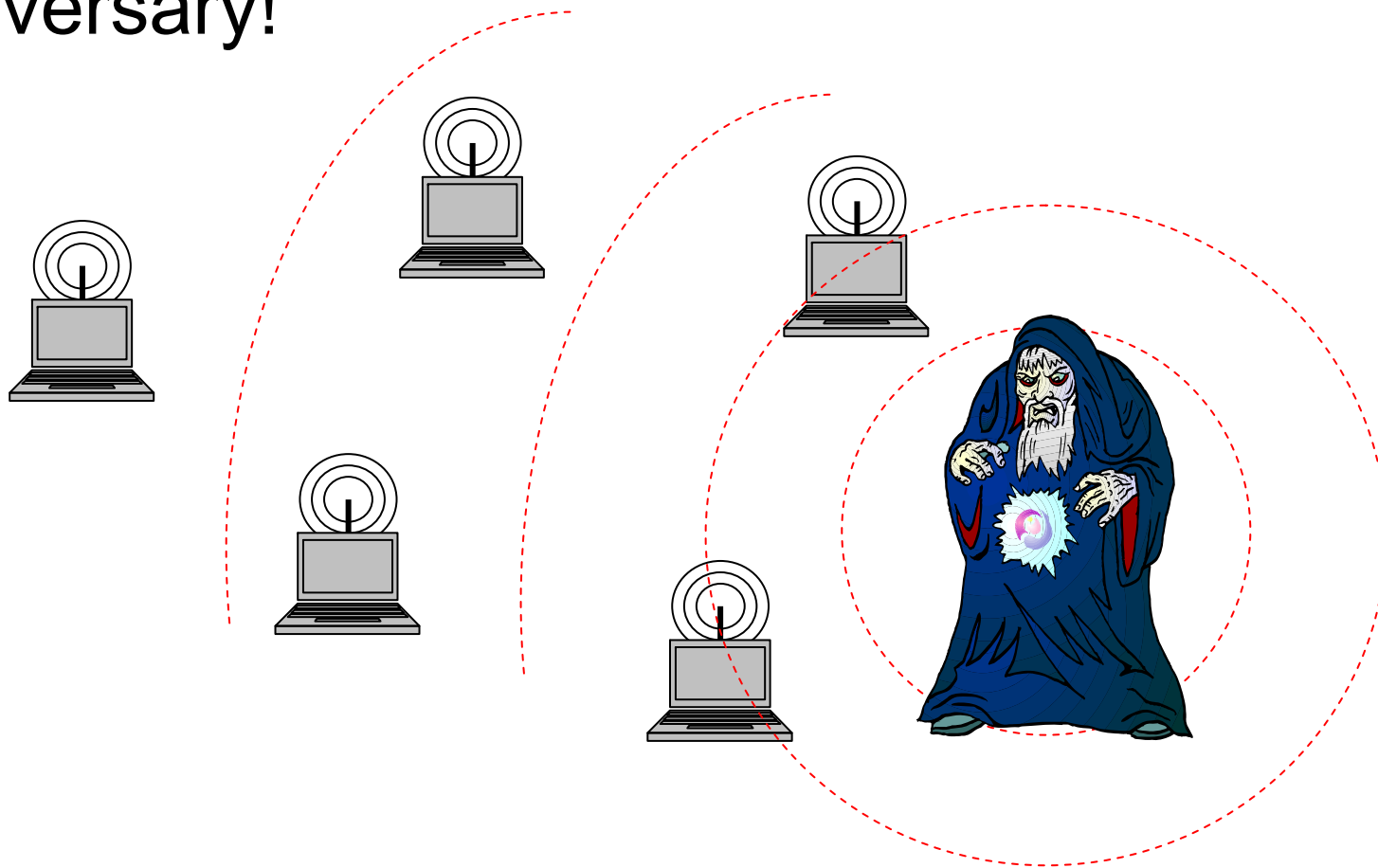
Our Approach: Adversarial Jamming

Idea: model unpredictable behaviors via adversary!



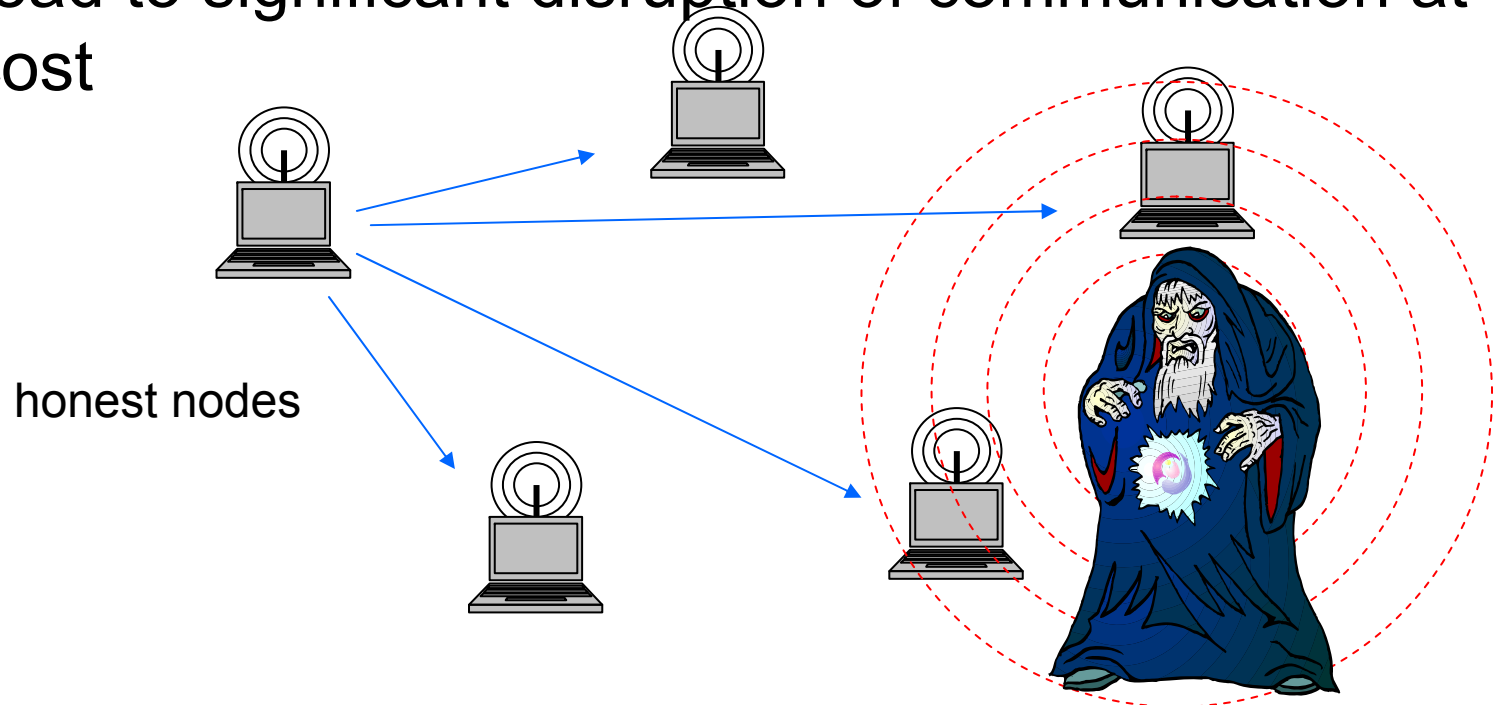
Our Approach: Adversarial Jamming

Idea: model unpredictable behaviors via adversary!



Adversarial physical layer jamming

- an adversary listens to the open medium and broadcasts in the same frequency band as the network
 - no special hardware required
 - can lead to significant disruption of communication at low cost



Adaptive adversary

- **Adaptive**: knows **protocol and entire history**
- The adversary is bounded **with respect to energy**, i.e., it has a certain budget **B** for each node **v**.
- **(B, T)-bounded adversary**: in any time window of size $w \geq T$, the adversary can add $\leq w * B / T$ to the noise level of each node **v**.

Signal-to-Interference Ratio

$$\frac{P_v(u)}{N + \sum_{w \in S} P_v(w) + J} \geq \beta$$

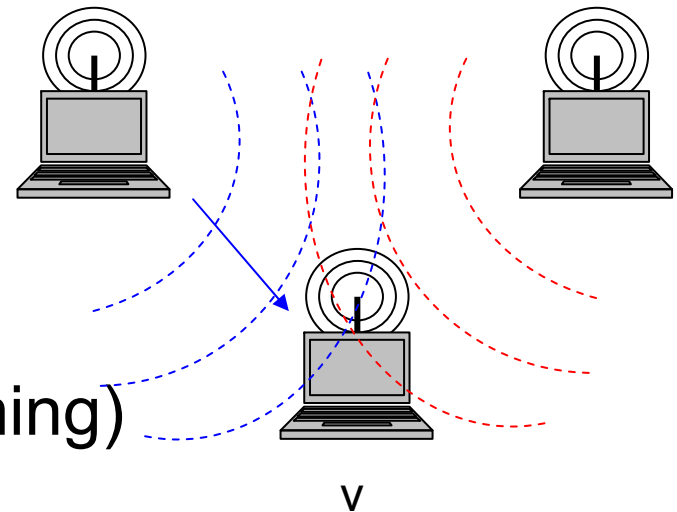
- Move away from any graph model.
- Combined signals from far away nodes can be strong enough to interrupt communications.
- Interference can come from anywhere.

Signal-to-Interference Ratio

- Define N_v as the overall noise level at node v .
- No clear distinctions between “idle” and “busy” anymore. A noise threshold τ_v is used by each node v to distinguish the two cases.
 - “idle”, when $N_v < \tau_v$.
 - “busy”, when $N_v > \tau_v$.
- τ_v is adjusted in an adaptive fashion, based on the events observed by node v

Wireless communication model

- at each time step, a node may decide to transmit a packet (nodes continuously contend to send packets)
- a node may transmit *or* sense the channel at any time step (half-duplex)
- when **sensing** the channel a node v may
 - **sense** an “idle” channel
 - **receive** a packet
 - **sense** a “busy” channel (cannot distinguish between message collisions and adversarial jamming)



SINR is difficult

- Define p as the cumulative transmission probability of all nodes.
- p is large, and τ_v is large:
 - idle? busy?
- p is small, and τ_v is small:
 - idle? busy?

Hard to tell.

Goal: Maximize the Throughput

- Define *Throughput* as the average number of messages successfully received by individual node in the network per round.

SINR-MAC Protocol

- each node v maintains
 - probability value p_v ,
 - noise level threshold τ_v
 - time window threshold T_v
 - counter c_v , and
 - $\gamma = O(1/(\log T + \log \log n))$
- Initially, $T_v = c_v = 1$, $\tau_v = 1$, and $p_v = p_{max} (< 1/24)$.
- synchronized time steps (for ease of explanation)

Basic approach

- a node v adapts p_v based only on steps when an “idle” channel or a successful message transmission are observed, ignoring all other steps (including all the blocked steps when the adversary transmits!)

time →



- $N_v < \tau_v$, “idle” steps
- successful transmissions
- $N_v > \tau_v$, “busy” steps

Basic approach

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time →



- $N_v < \tau_v$, “idle” steps
- successful transmissions
- $N_v > \tau_v$, “busy” steps

Basic approach

- a node v adapts τ_v based only on steps when an “idle” channel or a “busy” channel is observed!
 - τ_v is decreased when channel is “idle” .
 - τ_v is increased when channel is “busy”.

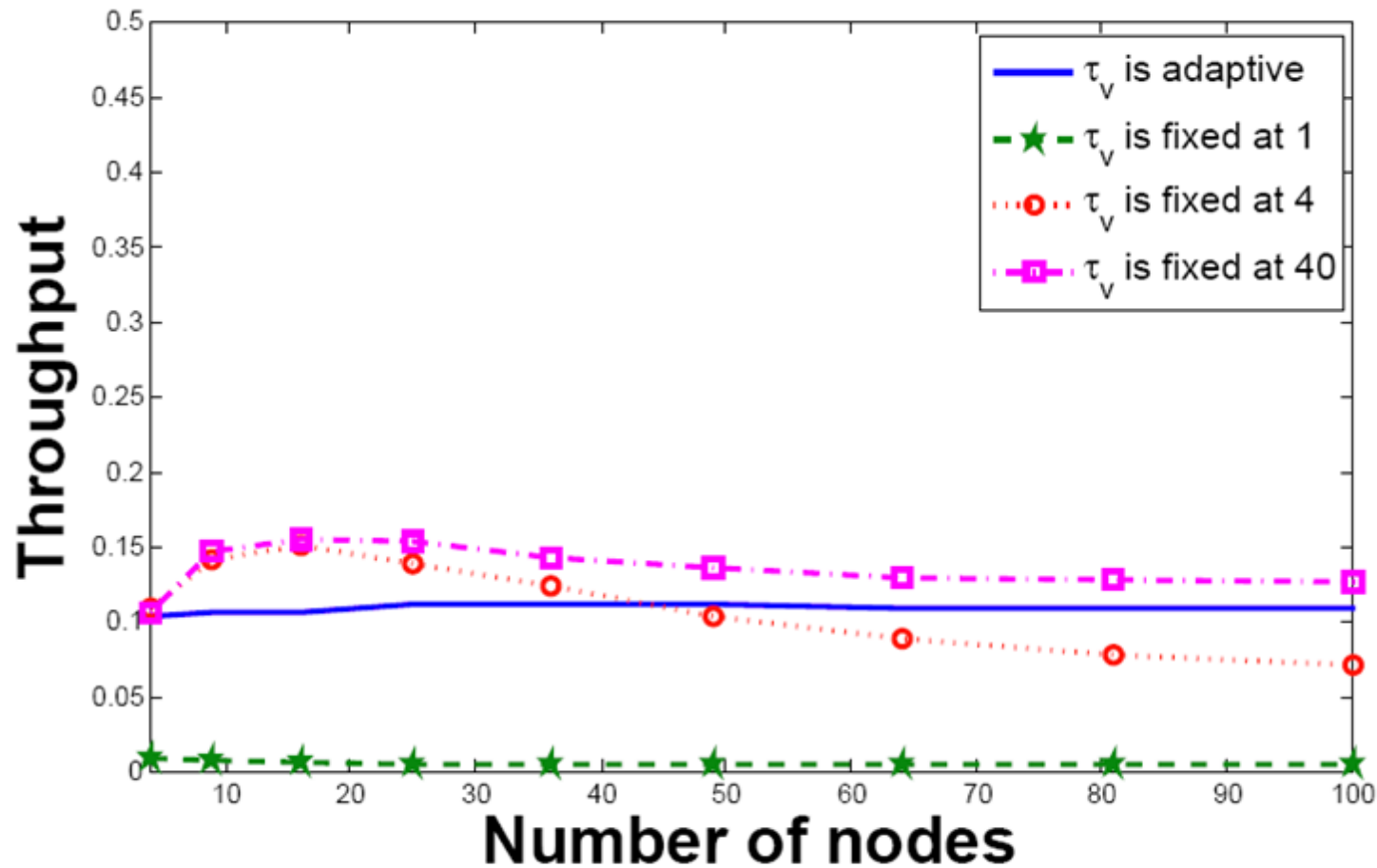
time →



- $N_v < \tau_v$, “idle” steps
- successful transmissions
- $N_v > \tau_v$, “busy” steps

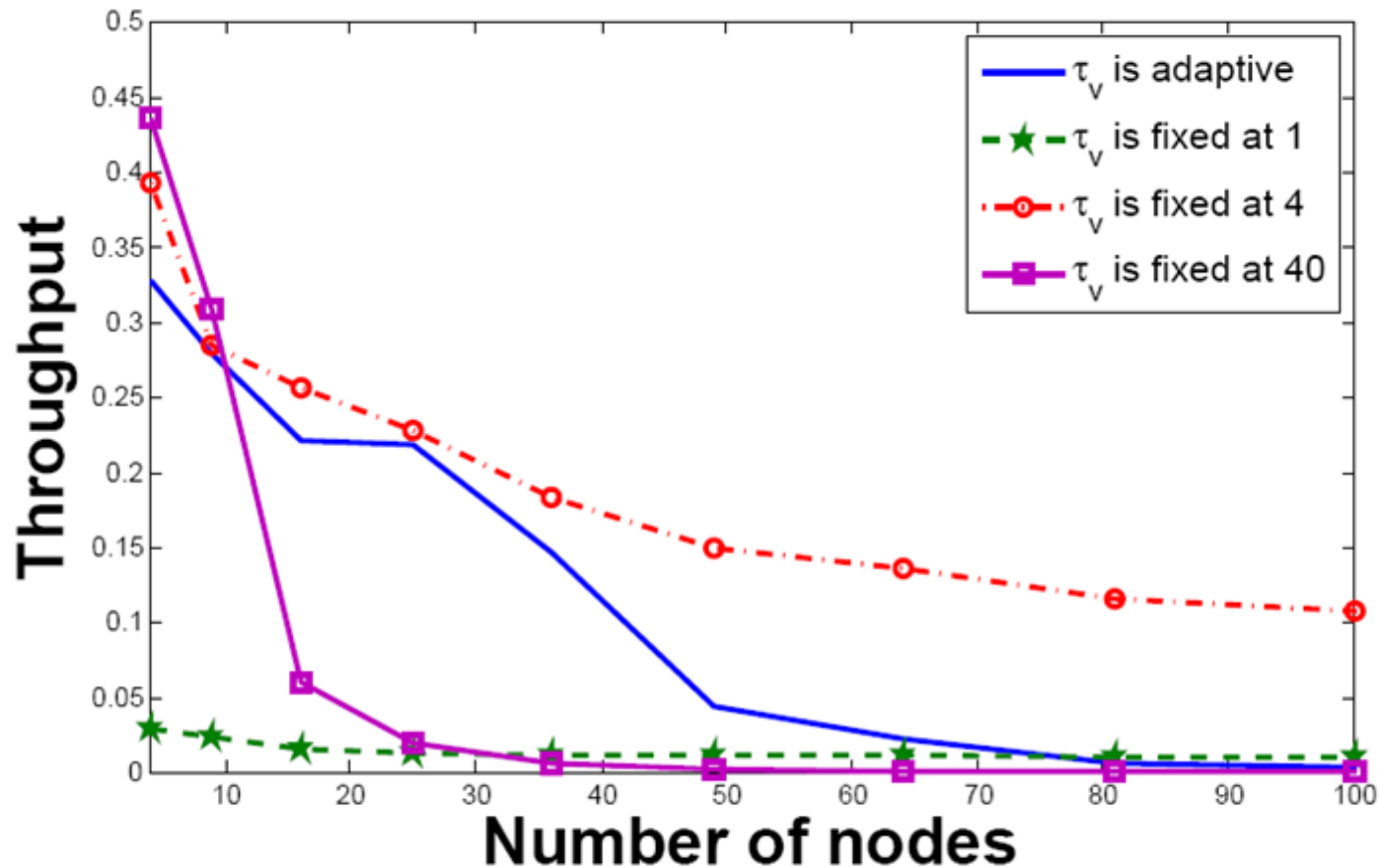
Simulation Result

Experiment 1: Throughput evaluation, $\rho_{max} = 1/24$.



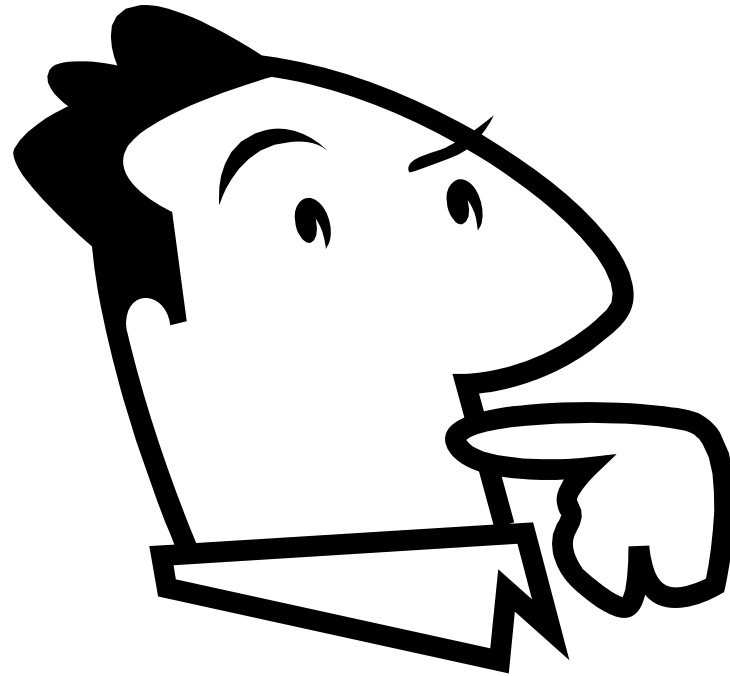
Simulation Result

Experiment 2: Throughput evaluation, $\rho_{max} = 1/2$.



Related Work

- Our jamming model:
 - Awerbuch, Richa, Scheideler (PODC 2008):
MAC protocol for single-hop network
 - Richa, Scheideler, Schmid, Zhang (DISC 2010): multi-hop networks
 - Richa, Scheideler, Schmid, Zhang (ICDCS 2011): reactive jammer
 - Richa, Scheideler, Schmid, Zhang (MOBIHOC 2011): leader election



Questions?